



Note

Colouring triangle-free intersection graphs of boxes  
on the planeA.V. Kostochka<sup>a,1</sup>, I.G. Perepelitsa<sup>b,\*2</sup><sup>a</sup>*Institute of Mathematics, 630090 Novosibirsk, Russia*<sup>b</sup>*Novosibirsk State University, 630090 Novosibirsk, Russia*

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**Abstract**

We prove that intersection graphs of boxes on the plane with girth 6 and 8 are 3- and 2-degenerate, respectively. This implies that these graphs are 4- and 3-list-colourable, respectively. © 2000 Published by Elsevier Science B.V. All rights reserved.

*Keywords:* Colouring; Intersection graphs; Maximum average degree

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**1. Introduction**

Among other interesting classes of intersection graphs, Asplund and Grünbaum [1] considered class  $\mathcal{B}$  of intersection graphs of *boxes*, i.e. rectangles with sides parallel to the axes on the plane. They proved that every graph in  $\mathcal{B}$  with clique number at most  $r$  is  $4r^2 - 3r$ -colourable. For  $r = 2$ , they obtained the exact result:

**Theorem 1** (Asplund and Grünbaum [1]). *The maximum chromatic number of a triangle-free graph in  $\mathcal{B}$  is six.*

Note that for three-dimensional boxes, the situation is different. Burling [2] constructed a series of triangle-free intersection graphs of three-dimensional boxes with arbitrarily high chromatic number.

The aim of the present paper is to show that graphs in  $\mathcal{B}$  with girth greater than 5 can be coloured with fewer than six colours.

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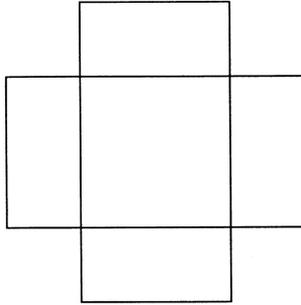


Fig. 1. An example of an  $x$ -intersection.

**Theorem 2.** *If a graph  $G \in \mathcal{B}$  has girth at least 6 or 8, then  $G$  is 4- or 3-list-colourable, respectively.*

In fact, a stronger statement will be proved. Recall that the *maximum average degree*,  $\text{mad}(G)$ , of a graph  $G$  is  $\max_{\emptyset \neq H \subseteq G} 2|E(H)|/|V(H)|$ , and a graph is said to be  $k$ -degenerate if every subgraph has a vertex of degree at most  $k$ . Since every graph  $G$  with  $\text{mad}(G) < k + 1$  is  $k$ -degenerate and every  $k$ -degenerate graph is  $(k + 1)$ -list-colourable, Theorem 2 is an immediate corollary of the following statement.

**Theorem 2'.** *If a graph  $G \in \mathcal{B}$  has girth at least 6 or 8, then  $\text{mad}(G)$  is strictly less than 4 or 3, respectively.*

Recall that ‘girth at least 4’ means ‘triangle-free’, and Theorem 1 applies. We do not know the situation with girth 5.

In the next section, we introduce the auxiliary notion of a *graphical graph* of a family of boxes and discuss some its properties. Theorem 2' is proved in the last section.

## 2. Graphical graphs

Let  $\mathcal{A}$  be a family of boxes on the plane whose intersection graph is triangle-free. We shall say that two boxes  $B_1$  and  $B_2$  *x-intersect* if their intersection is nonempty but no corner of  $B_i$  is covered by  $B_{3-i}$ ,  $i = 1, 2$  (see Fig. 1).

We also say that  $B_1$  and  $B_2$  *z-intersect* (respectively, *y-intersect*) if exactly one corner of  $B_i$  is covered by  $B_{3-i}$ ,  $i = 1, 2$  (respectively, one of them contains exactly two corners of the other) (see Fig. 2).

For every  $B \in \mathcal{A}$ , let  $d_{zy}(B)$  denote the number of boxes  $B' \in \mathcal{A}$  such that  $B$  and  $B'$   $z$ - or  $y$ -intersect, and let  $d_x(B)$  denote the number of boxes  $B'' \in \mathcal{A}$  such that  $B$  and  $B''$   $x$ -intersect.

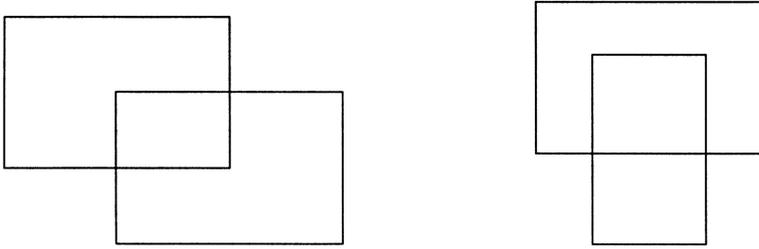


Fig. 2. Examples of a z-intersection (left) and a y-intersection (right).

Let  $G \in \mathcal{B}$ . There are infinitely many families  $\mathcal{A}$  of boxes on the plane such that  $G$  is the intersection graph of  $\mathcal{A}$ . We say that some such  $\mathcal{A}$  is a *standard representation* of  $G$  if all the sides of all boxes lay on different lines and the sum  $\sum_{B \in \mathcal{A}} d_x(B)$  is minimum possible.

**Lemma 1.** *Let a triangle-free  $G \in \mathcal{B}$  have minimum degree at least 2. Let  $\mathcal{A}$  be a standard representation of  $G$ . Then  $d_{zy}(B) \geq 2$  for every  $B \in \mathcal{A}$ .*

**Proof.** Let  $B \in \mathcal{A}$ . If  $d_x(B) = 0$ , then we are done, since  $d_x(B) + d_{zy}(B) \geq 2$ . Suppose that  $B$  and  $B'$  x-intersect. Then  $B \setminus B'$  is the disjoint union of two boxes, say,  $B_1$  and  $B_2$ . Since  $G$  is triangle-free, no box in  $\mathcal{A} - B$  meets both  $B_1$  and  $B_2$ . If  $B_1$  does not meet other boxes, then instead of  $B$  we can use a shorter box  $\hat{B}$  which y-intersects  $B'$ . This contradicts the condition that  $\mathcal{A}$  is a standard representation of  $G$ . It follows that each of  $B_1$  and  $B_2$  meets another box, say  $B'_1$  and  $B'_2$ , in  $\mathcal{A}$ . Repeating (if needed) the argument, we obtain finally that each of  $B_1$  and  $B_2$  z- or y-intersects some box in  $\mathcal{A}$ .  $\square$

The following notion (similar to one used in [3]) will be helpful. The vertices of the *graphical graph*,  $\tilde{G}(\mathcal{A})$ , of a family  $\mathcal{A}$  of boxes on the plane are the corners of all boxes and the crossing points of the sides of the boxes; the edges of  $\tilde{G}(\mathcal{A})$  are the pieces of sides of boxes connecting these points (see Fig. 3).

By the construction,  $\tilde{G}(\mathcal{A})$  is always a plane graph. We say that a face of  $\tilde{G}(\mathcal{A})$  is of type  $i$ , if its interior is covered by exactly  $i$  boxes of  $\mathcal{A}$ . The set of faces of type  $i$  will be denoted by  $T_i = T_i(\mathcal{A})$ . The faces of type 0 will be called *outer* and the other faces will be called *interior*. For every face  $F$  of  $\tilde{G}$ , let  $r(F)$  denote its *rank*, i.e. the number of edges on the boundary of  $F$ . Observe that (since the intersection graph of  $\mathcal{A}$  is triangle-free) there are no faces of type  $i$  for  $i \geq 3$  and that the faces of type 2 are always of rank four.

Let  $G \in \mathcal{B}$  be triangle-free,  $\mathcal{A}$  be a standard representation of  $G$  and  $\tilde{G} = \tilde{G}(\mathcal{A})$  be the graphical graph of  $\mathcal{A}$ . Let  $n = n(G)$  and  $m = m(G)$  denote the number of vertices and edges of  $G$ , respectively. Furthermore, let  $v = v(\tilde{G})$ ,  $e = e(\tilde{G})$ ,  $f = f(\tilde{G})$  and  $f_i = f_i(\tilde{G})$  denote the number of vertices, edges, faces and faces of rank  $i$  in  $\tilde{G}$ ,

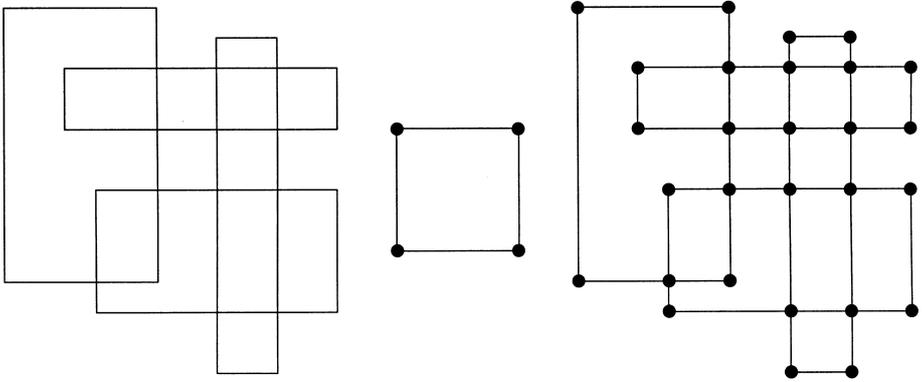


Fig. 3. An example: a family of boxes (left), its intersection graph (in the center) and its graphical graph (right).

respectively. Then the following equalities hold:

$$2m = \sum_{B \in \mathcal{A}} d_x(B) + \sum_{B \in \mathcal{A}} d_{zy}(B), \tag{1}$$

$$v = 4n + 2 \sum_{B \in \mathcal{A}} d_x(B) + \sum_{B \in \mathcal{A}} d_{zy}(B), \tag{2}$$

$$e = 4n + 4 \sum_{B \in \mathcal{A}} d_x(B) + 2 \sum_{B \in \mathcal{A}} d_{zy}(B), \tag{3}$$

$$2e = \sum_{i=4}^{\infty} i \cdot f_i. \tag{4}$$

If a box participates in  $j$   $x$ -intersections, then it contains exactly  $j + 1$  faces of type 1 in  $\tilde{G}$ . Hence

$$|T_1| = n + \sum_{B \in \mathcal{A}} d_x(B). \tag{5}$$

Each  $z$ -intersection creates a new face of type 2 and increases the sizes of both faces of type 1 adjacent to this new face by two. Each  $y$ -intersection creates a new face of type 2 and increases the size of one of the faces of type 1 adjacent to this new face by four. This yields the following equation:

$$\sum_{F \in T_1} r(F) = 4|T_1| + 2 \sum_{B \in \mathcal{A}} d_{zy}(B). \tag{6}$$

### 3. Proof of Theorem 2'

Suppose that the theorem does not hold. Let  $g \in \{6, 8\}$  and  $G$  be a counter-example to the theorem for this  $g$  with the smallest number of vertices. Let  $\mathcal{A}$  be a standard representation of  $G$  and  $\tilde{G} = \tilde{G}(\mathcal{A})$  be the graphical graph of  $\mathcal{A}$ .

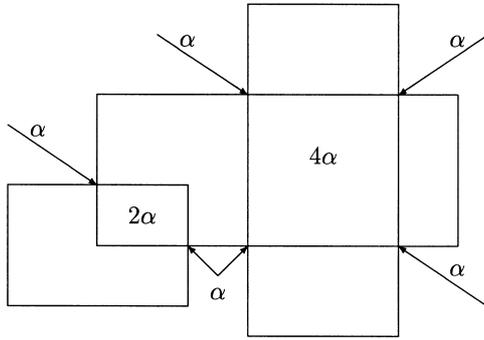


Fig. 4.

Due to the minimality of  $G$ , we have  $2|E(G)|/|V(G)| \geq 4$  if  $g = 6$  and  $2|E(G)|/|V(G)| \geq 3$  if  $g = 8$ . If  $G$  has a vertex  $B$  of degree at most one, then  $G - B$  is also a counter-example to the theorem, a contradiction. Thus, the minimum degree of  $G$  is at least two. Hence Lemma 1 and Eqs. (1)–(6) hold for  $G$ ,  $\mathcal{A}$  and  $\tilde{G}$ . Let  $T_2^x$  denote the set of faces in  $T_2$  created by  $x$ -intersections, and  $T_2^{zy} = T_2 \setminus T_2^x$ . In order to estimate  $f = \sum_{F \in T_0 \cup T_1 \cup T_2} r(F)$  we use discharging, i.e. giving each  $F$  a number  $r'(F)$  following certain rules so that  $\sum_{F \in T_0 \cup T_1 \cup T_2} r'(F) = \sum_{F \in T_0 \cup T_1 \cup T_2} r(F)$ .

*$\alpha$ -Discharging procedure:* For each  $F \in T_1$ , set  $r'(F) = r(F)$ . For each  $F \in T_2^x$ , set  $r'(F) = r(F) + 4\alpha$ . For each  $F \in T_2^{zy}$ , set  $r'(F) = r(F) + 2\alpha$ . For each  $F \in T_0$ , set  $r'(F) = r(F) - a(F, \pi/2)\alpha$ , where  $a(F, \phi)$  denotes the number of angles equal to  $\phi$  on the boundary of  $F$ .

Since each angle of  $\pi/2$  on the boundary of a face  $F \in T_0$  is vertical to an angle on the boundary of a face of type 2, we can consider our procedure as if every  $F \in T_0$  sends through each angle equal to  $\pi/2$  on its boundary the share  $\alpha$  to the face of type 2 incident with the vertex of this angle (see Fig. 4).

This proves that  $\sum_{F \in T_0 \cup T_1 \cup T_2} r'(F) = \sum_{F \in T_0 \cup T_1 \cup T_2} r(F)$ .

**Lemma 2.** *If*

$$\alpha \leq \frac{6g - 24}{4 + 3g}, \tag{7}$$

then for every  $F \in T_0$ ,  $r'(F) \geq 4 + 4\alpha/3$ .

**Proof.** Let  $F \in T_0$ . Condition (7) is equivalent to the inequality

$$(2 - \alpha)g - 4 \geq 4 + \frac{4}{3}\alpha.$$

Note that the number of boxes adjacent to  $F$  is at most  $a(F, \pi/2)$ . Therefore,  $a(F, \pi/2) \geq g$  and

$$(2 - \alpha)a(F, \pi/2) - 4 \geq 4 + \frac{4}{3}\alpha.$$

In other words,

$$(2a(F, \pi/2) - 4) - \alpha a(F, \pi/2) \geq 4 + \frac{4}{3}\alpha. \tag{8}$$

Observe that for every face  $F$  whose angles are only  $\pi/2$  and  $3\pi/2$ ,  $a(F, \pi/2) - a(F, 3\pi/2) = 4$ , that is,

$$r(F) = 2a(F, \pi/2) - 4.$$

This together with (8) yields that  $r(F) - \alpha a(F, \pi/2) \geq 4 + \frac{4}{3}\alpha$  which proves the lemma.  $\square$

Let (7) hold. Then due to the discharging procedure and Lemma 2,

$$2e = \sum_{F \in T_0 \cup T_1 \cup T_2} r'(F) \geq \left(4 + \frac{4\alpha}{3}\right) |T_0| + \sum_{F \in T_1} r(F) + (4 + 2\alpha) |T_2| + 2\alpha |T_2^x|.$$

By (6) and (5),

$$\begin{aligned} \sum_{F \in T_1} r(F) &= 4|T_1| + 2 \sum_{B \in \mathcal{B}} d_{zy}(B) \\ &= \left(4 + \frac{4\alpha}{3}\right) |T_1| - \frac{4\alpha}{3} \left(n + \sum_{B \in \mathcal{B}} d_x(B)\right) + 2 \sum_{B \in \mathcal{B}} d_{zy}(B). \end{aligned}$$

Therefore,

$$2e \geq \left(4 + \frac{4\alpha}{3}\right) f - \frac{4\alpha}{3} \left(n + \sum_{B \in \mathcal{B}} d_x(B)\right) + 2 \sum_{B \in \mathcal{B}} d_{zy}(B) + \frac{2\alpha}{3} |T_2^{zy}| + \frac{8\alpha}{3} |T_2^x|.$$

Since  $|T_2^x| = \frac{1}{2} \sum_{B \in \mathcal{B}} d_x(B)$  and  $|T_2^{zy}| = \frac{1}{2} \sum_{B \in \mathcal{B}} d_{zy}(B)$ , this is equivalent to

$$2e \geq \left(4 + \frac{4\alpha}{3}\right) f - \frac{4\alpha}{3} n + \left(2 + \frac{\alpha}{3}\right) \sum_{B \in \mathcal{B}} d_{zy}(B).$$

Recall that  $v - e + f = 2$  according to Euler’s formula. Hence

$$2e > \left(4 + \frac{4\alpha}{3}\right) (e - v) - \frac{4\alpha}{3} n + \left(2 + \frac{\alpha}{3}\right) \sum_{B \in \mathcal{B}} d_{zy}(B).$$

Substituting the expressions for  $v$  and  $e$  from (2) and (3), we get

$$\begin{aligned} 8n + 8 \sum_{B \in \mathcal{B}} d_x(B) + 4 \sum_{B \in \mathcal{B}} d_{zy}(B) \\ > \left(4 + \frac{4\alpha}{3}\right) \left(2 \sum_{B \in \mathcal{B}} d_x(B) + \sum_{B \in \mathcal{B}} d_{zy}(B)\right) - \frac{4\alpha}{3} n + \left(2 + \frac{\alpha}{3}\right) \sum_{B \in \mathcal{B}} d_{zy}(B). \end{aligned}$$

In other words,

$$\begin{aligned} \left(8 + \frac{4\alpha}{3}\right) n &> \frac{8\alpha}{3} \sum_{B \in \mathcal{B}} d_x(B) + \left(2 + \frac{5\alpha}{3}\right) \sum_{B \in \mathcal{B}} d_{zy}(B) \\ &= \frac{16\alpha}{3} m + (2 - \alpha) \sum_{B \in \mathcal{B}} d_{zy}(B). \end{aligned}$$

By Lemma 1,  $\sum_{B \in \mathcal{A}} d_{zy}(B) \geq 2n$  and hence

$$\left(8 + \frac{4\alpha}{3}\right)n > \frac{16\alpha}{3}m + (2 - \alpha)2n,$$

i.e.

$$\frac{m}{n} < \frac{6 + 5\alpha}{8\alpha} = \frac{3}{4\alpha} + \frac{5}{8}. \quad (9)$$

Let  $g \geq 8$ . Then  $\alpha = \frac{6}{7}$  satisfies (7), and from (9) we derive that

$$\frac{m}{n} < \frac{7}{8} + \frac{5}{8} = \frac{3}{2}.$$

Similarly, for  $g \geq 6$ ,  $\alpha = \frac{6}{11}$  satisfies (7), and

$$\frac{m}{n} < \frac{11}{8} + \frac{5}{8} = 2.$$

A contradiction to the choice of  $G$  proves Theorem 2'.

**Remark.** Maybe the chromatic number of intersection graphs of three-dimensional boxes with large girth is bounded.

## Acknowledgements

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