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# On equitable $\triangle$ -coloring of graphs with low average degree $\stackrel{\text{tr}}{\leftarrow}$

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### 7 Abstract

An *equitable coloring* of a graph is a proper vertex coloring such that the sizes of any two color classes differ by at most 1. Hajnal and Szemerédi proved that every graph with maximum degree  $\Delta$  is equitably *k*-colorable for every  $k \ge \Delta + 1$ . Chen, Lih, and Wu conjectured that every connected graph with maximum degree  $\Delta \ge 3$  distinct from  $K_{\Delta+1}$  and  $K_{\Delta,\Delta}$  is equitably  $\Delta$ -colorable. This conjecture has been proved for graphs in some classes such as bipartite graphs, outerplanar graphs, graphs with maximum degree

- 3, interval graphs. We prove that this conjecture holds for graphs with average degree at most  $\Delta/5$ .
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#### 15 1. Introduction

In several applications of coloring as a partition problem there is an additional requirement that color classes be not so large or be of approximately the same size. Examples are the mutual exclusion scheduling problem [1,17], scheduling in communication systems [7], construction timetables [9], and round-a-clock scheduling [18]. For other

- 19 applications in scheduling, partitioning, and load balancing problems, one can look into [2,12,17]. A model imposing such a requirement is *equitable coloring*—a proper coloring such that color classes differ in size by at most one. A
- 21 good survey on equitable colorings of graphs is given in [13]. Recently, Pemmaraju [16] and Janson and Ruciński [8] used equitable colorings to derive deviation bounds for sums of dependent random variables that exhibit limited

dependence.Unlike in the case of ordinary coloring, a graph may have an equitable *k*-coloring (i.e., an equitable coloring with *k* 

colors) but have no equitable (k + 1)-coloring. For example, the complete bipartite graph  $K_{2n+1,2n+1}$  has the obvious equitable 2-coloring, but has no equitable (2n + 1)-coloring. Thus, it is natural to look for the minimum number, eq(G),

such that for every  $k \ge eq(G)$ , *G* has an equitable *k*-coloring.

The difficulty of finding eq(G) is not less than that of finding the chromatic number. Thus, already in the class of planar graphs, finding eq(G) is an NP-hard problem. This situation prompted studying extremal problems on relations of

eq(*G*) with other graph parameters. Hajnal and Szemerédi [6] settled a conjecture of Erdős by proving that eq(*G*)  $\leq \Delta + 1$ for every graph *G* with maximum degree at most  $\Delta$ . This bound is sharp, as shows the example of  $K_{2n+1,2n+1}$  above.

Other natural examples showing sharpness of Hajnal–Szemerédi Theorem are graphs with chromatic number greater

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1 than maximum degree, i.e. complete graphs and odd cycles. Chen et al. [3] proposed the following analogue of Brooks' Theorem for equitable coloring:

3 **Conjecture 1** (*Chen et al.* [3]). Let G be a connected graph with maximum degree  $\Delta$ . If G is distinct from  $K_{\Delta+1}, K_{\Delta,\Delta}$ , and is not an odd cycle, then G has an equitable coloring with  $\Delta$  colors.

They proved the conjecture for graphs with maximum degree at most three. Later, Yap and Zhang [20,21] proved that the conjecture holds for outerplanar graphs and planar graphs with maximum degree at least 13. Nakprasit (unpublished)
extended the result of Yap and Zhang [21] to planar graphs with maximum degree at least 9. Lih and Wu [14] verified the conjecture for bipartite graphs, and Chen et al. [4] verified it for interval graphs. Kostochka et al. [11] studied a list

- 9 analogue of equitable coloring and proved the validity of Conjecture 1 for equitable list coloring in classes of interval graphs and 2-degenerate graphs. It follows from [10] that the conjecture holds for *d*-degenerate graphs with maximum
- degree *∆* if *d* ≤ (*∆* − 1)/14.
  Recall that a graph *G* is *d*-degenerate, if each subgraph *G'* of *G* has a vertex of degree (in *G'*) at most *d* (see, e.g., [19,
  p. 269]). In other words, one can destroy any *d*-degenerate graph by successively deleting vertices of degree at most
- *d*. Forests are exactly 1-degenerate graphs. It is also well known that every outerplanar graph is 2-degenerate (see,
- e.g., [19, p. 240]), and every planar graph is 5-degenerate. To say that a graph has 'low degeneracy' is about the same as to say that every subgraph of *G* has a 'small average degree'. In this paper, we prove Conjecture 1 for graphs that
- 17 have 'low average degree' themselves without restrictions on average degrees of subgraphs.

**Theorem 1.** Let  $\Delta$ ,  $n \ge 46$ . Suppose that an n-vertex graph G = (V, E) has maximum degree at most  $\Delta$  and  $|E| \le \Delta n/5$ . 19 If  $K_{\Delta+1}$  is not a subgraph of G, then G has an equitable coloring with  $\Delta$  colors.

An immediate consequence of Theorem 1 is that Conjecture 1 holds for *d*-degenerate graphs with maximum degree  $\Delta$  if  $d \leq \Delta/10$ .

In order to prove Theorem 1, we will need the following statement.

**Theorem 2.** Let  $\Delta \ge 3$  and G be a  $K_{\Delta+1}$ -free graph with  $\Delta(G) \le \Delta$ . Suppose that G - v has a  $\Delta$ -coloring with color classes  $M_1, M_2, \ldots, M_{\Delta}$ . Then G has a  $\Delta$ -coloring with color classes  $M'_1, M'_2, \ldots, M'_{\Delta}$  such that  $|M_i| = |M'_i|$  for all *i* apart from one.

We call this statement a theorem, since it seems that it has its own merit. It can be considered as a slight refinement of the Brooks' Theorem. For example, Theorem 2 has the following easy consequence.

**Corollary 1.** Let  $\Delta \ge 3$  and G be a  $K_{\Delta+1}$ -free graph with  $\Delta(G) \le \Delta$ . Let  $1 \le k \le \Delta$  and  $m_k$  be the order of a maximum k-colorable subgraph of G. Then there is a proper coloring of G with at most  $\Delta$  colors in which the total number of vertices in some k color classes is  $m_k$ . In particular, G has a proper coloring with at most  $\Delta$  colors in which one of the color classes has  $\alpha(G)$  vertices.

The first step of the proof of Theorem 1 uses the Hajnal–Szemerédi Theorem whose complexity we have not analyzed. The rest of the proof can be rewritten as a polynomial time algorithm for equitable coloring of a graph.

The structure of the paper is as follows. In the next section, we prove Theorem 2. In Section 3 we introduce some useful notions and verify Theorem 1 for graphs with few vertices. In Section 4, we prove the bounded-size version of

Theorem 1: we demand each color class to be of size at most  $\lceil n/\Delta \rceil$  but allow 'small' color classes. We finalize the proof of Theorem 1 in the last section.

Throughout the paper, we use standard graph-theoretic definitions and notation (see, e.g., [5,19]).

### 39 2. Proof of Theorem 2

The proof follows the steps of a proof of Brooks' Theorem by Mel'nikov and Vizing [15] (see also [5, p. 99]).

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- Let Δ≥3 and G be a K<sub>Δ+1</sub>-free graph with Δ(G) ≤ Δ. Let f be a proper coloring of G v with at most Δ colors. Suppose that M<sub>1</sub>, M<sub>2</sub>,..., M<sub>Δ</sub> are color classes (maybe empty) of f and that G has no Δ-coloring with color classes
   M'<sub>1</sub>, M'<sub>2</sub>,..., M'<sub>Δ</sub> such that |M<sub>i</sub>| = |M'<sub>i</sub>| for all i apart from one. Let G<sub>ij</sub> denote the subgraph of G induced by M<sub>i</sub> ∪ M<sub>j</sub> ∪ {v}, and G'<sub>ij</sub> denote the component of G<sub>ij</sub> containing v. We will deliver the proof in a series of claims.
- 5 **Claim 2.1.** The vertex v has exactly one neighbor in each  $M_i$ .

**Proof.** Otherwise, there is a color class  $M_l$  containing no neighbors of v. We simply add v to  $M_l$ . Since other color classes do not change, the conclusion of the theorem holds, a contradiction.  $\Box$ 

In view of this claim, let  $w_i$  denote the only neighbor of v in  $M_i$ .

9 **Claim 2.2.** For each  $1 \le i < j \le \Delta$ , the graph  $G'_{ij}$  is an odd cycle.

**Proof.** Suppose that some  $G'_{ij}$  is not an odd cycle. By Claim 2.1,  $G'_{ij}$  cannot be an even cycle. Thus  $G'_{ij}$  has a

11 vertex of degree (in  $G'_{ij}$ ) distinct from 2. Let u be a vertex closest to v in  $G'_{ij}$  with degree in  $G'_{ij}$  not equal to 2. Let  $P = (v, v_1, \ldots, v_k)$ , where  $v_k = u$ , be the shortest v, u-path in  $G'_{ij}$  (since  $G'_{ij} - v$  is bipartite, this path is unique).

- 13 Note that each internal vertex of P has degree 2 in  $G'_{ij}$ . Denote  $A_j = \{v_1, \dots, v_k\} M_i$ , and  $A_i = \{v_1, \dots, v_k\} M_j$ . Let  $B_i = M_i \cup A_j - A_i$ , and  $B_j = M_j \cup A_i - A_j$ . We may assume that  $v_1 = w_i$ .
- 15 Case 1:  $d_{G'_{ij}}(u) = 1$ . Then P v is a component in  $G_{ij} v$ . Therefore, the sets  $B_i$  and  $B_j$  are independent. Note that  $B_i$  does not contain neighbors of v, since the only neighbor,  $v_1$ , of v in  $M_i$  belongs to  $B_j$  and  $A_j$  does not contain
- any neighbor of v. Thus  $M'_i = B_i \cup \{v\}$ ,  $M'_j = B_j$  and  $M'_m = M_m$ , for  $m \neq i$ , j, are color classes of a proper coloring of G. If k is odd, then  $|M'_i| = |M_i|$ , and if k is even, then  $|M'_i| = |M_j|$ . This contradicts the choice of G and f.
- 19 *Case* 2:  $d_{G'_{ij}}(u) \ge 3$ . Since  $d_G(u) \le \Delta$  and  $d_{G_{ij}}(u) \ge 3$ , there is a color class  $M_l$ ,  $l \ne i, j$ , with no neighbors of u. Hence  $M_l \cup \{u\}$  is independent. Similarly to Case 1, the sets  $B_i - v_k$  and  $B_j - v_k$  are independent. Thus,  $M'_i = B_i \cup \{v\}$ ,
- 21  $M'_j = B_j, M'_l = M_l \cup \{u\}$  and  $M'_m = M_m$ , for  $m \neq i, j, l$ , are color classes of a proper coloring of G. Note that independently of the parity of k,  $|M'_j| = |M_i|$ , and  $|M'_j| = |M_j|$ . This contradicts the choice of G and f.  $\Box$
- 23 **Claim 2.3.** For any distinct *i*, *j* and *s*, the components  $G'_{ij}$  and  $G'_{js}$  share exactly two vertices, namely, the vertex *v* and the neighbor,  $w_i$ , of *v* in  $M_j$ .
- 25 **Proof.** By the definition,  $\{v, w_j\} \subseteq G'_{ij} \cap G'_{js}$ . Suppose  $u \in G'_{ij} \cap G'_{js} \{v, w_j\}$ . By Claim 2.2, *u* has four neighbors in  $G'_{ij} \cup G'_{is}$ . Hence there is a color class  $M_l$ ,  $l \neq j$ , with no neighbors of *u*. Let  $v, v_1, \ldots, v_k, u$  be the *v*, *u*-path
- 27 in  $G'_{ij}$  with  $v_1 = w_i$ . Note that k is odd, since  $u \in M_j$  and  $v_1 = w_i \in M_i$ . Denote  $A_j = \{v_1, \ldots, v_k\} M_i$ , and  $A_i = \{v_1, \ldots, v_k\} M_j$ . Let  $B_i = M_i \cup A_j A_i$ , and  $B_j = M_j \cup A_i A_j \{u\}$ . Similarly to the proof in Claim 2.2,
- 29 the sets  $B_i$  and  $B_j$  are independent. Also similarly to the proof of Claim 2.2, the sets  $M'_i = B_i \cup \{v\}$ ,  $M'_j = B_j$ ,  $M'_i = M_l \cup \{u\}$  and  $M'_m = M_m$ , for  $m \neq i, j, l$ , are color classes of a proper coloring of G. Moreover,  $|M_m| = |M'_m|$
- 31 for all *m* apart from m = l. This proves the claim.

Now, we are ready to finish the proof of Theorem 2. Among all colorings of G - v with color class sizes 33  $|M_1|, \ldots, |M_{\Delta}|$ , choose a coloring  $f = (M_1, \ldots, M_{\Delta})$  and indices *i* and *j* so that  $G'_{ij} - v = G'_{ij}(f) - v$  has the

- largest order. Recall that by the claims above,  $G'_{ij} v$  is a path with end-vertices  $w_i$  and  $w_j$ . We may assume that 35 i = 1, j = 2, and  $G'_{12} - v$  is a path  $P_1 = (v_1, \dots, v_q)$ , where  $v_1 = w_1$  and  $v_q = w_2$ . If q = 2, then by the maximality of  $G'_{ij}$ , all  $w_i$ s are adjacent to each other and together with v form a  $K_{\Delta+1}$ , a contradiction. Thus,  $q \ge 4$  and hence 37  $v_{q-1} \ne w_1$ .
- Similarly,  $G'_{23} v$  is a path,  $P_2$ , with end-vertices  $w_2$  and  $w_3$ . Let coloring  $f_1$  be obtained from f by swapping the colors 2 and 3 on vertices in  $G'_{23} v$ . Since  $G'_{23} v$  is a  $w_2$ ,  $w_3$ -path, the color class sizes in  $f_1$  are the same as in f.
- By Claim 2.2, the new bicolored subgraph  $H = G'_{12}(f_1) v$  has to be a  $w_1$ ,  $w_3$ -path (because  $w_3$  is the only neighbor of v in the second color class of  $f_1$ ). Note that H contains  $P_1 - v_a$ . By the maximality of q, the vertices  $v_{q-1}$  and  $w_3$
- 41 of v in the second color class of  $f_1$ ). Note that H contains  $P_1 v_q$ . By the maximality of q, the vertices  $v_{q-1}$  and  $w_3$  must be adjacent. But then  $v_{q-1}$  belongs to  $G'_{13}(f)$ , a contradiction to Claim 2.3. This proves the theorem.  $\Box$

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- 1 **Proof of Corollary 1.** Suppose that *H* is a *k*-colorable subgraph of *G* of the maximum order. Let  $V(G) V(H) = \{v_1, \ldots, v_i\}$ . Let  $G_0 = H$  and  $G_i = G[V(H) \cup \{v_1, \ldots, v_i\}]$  for  $i = 1, \ldots, t$ . Let  $f_0 = (M_1, \ldots, M_d)$  be a coloring
- 3 of *H* with  $M_{k+1} = M_{k+2} = \cdots = M_{\Delta} = \emptyset$ . Now, for i = 1, ..., t, apply Theorem 2 to  $G_i$  and  $f_{i-1}$  to produce a coloring  $f_i$  of  $G_i$ . Then by this theorem, the total number of vertices of colors 1, ..., k in  $G_t = G$  is at least |V(H)|,
- 5 but it cannot be greater by the maximality of *H*. This proves the corollary. As an additional feature, we have that none of the sizes of the first *k* color classes changed.  $\Box$
- 7 **Corollary 2.** Let  $\Delta \ge 3$  and G = (V, E) be a  $K_{\Delta+1}$ -free graph with  $\Delta(G) \le \Delta$ . Let  $|V| = n = k(\Delta + 1) + r$ , where  $0 \le r \le \Delta$ . Then G has a  $\Delta$ -coloring f with color classes  $M_1, M_2, \ldots, M_\Delta$  such that
  - (i)  $|M_i| \ge k$  for every *i*;

(ii) for every set Z of color classes,  $|\bigcup_{M \in \mathbb{Z}} M| \leq k + |Z| \lceil n/\Delta + 1 \rceil$ ; in particular,  $|M_i| \leq k + \lceil n/\Delta + 1 \rceil$  for every i; (iii) if  $|M_i| = k + 1 + p$  for some  $p \geq 1$ , then the degree of every  $v \in M_i$  in G is at least  $\Delta - (k + r - 1)/p$ .

Proof. By the Hajnal–Szemerédi Theorem, G has an equitable (∆ + 1)-coloring f'. Under the conditions of the corollary, exactly r color classes of f' have size k + 1. Let M' be a color class of f' with |M'| = k. Adding the vertices of M' one by one to G - M' and applying Theorem 2 on every step, we get a Δ-coloring f" of G satisfying (i) and (ii).

- Now, consider the following procedure: If a vertex v in a color class  $M_i$  of size z has no neighbors in a color class  $M_i$  of size at most z 2, then move v from  $M_i$  to  $M_i$ . Clearly, we will stop after a finite number of steps. We claim
- that the final  $\Delta$ -coloring f is what we need. Indeed, once a coloring satisfies (i) and (ii), such moves do not destroy these properties. Since we have stopped our procedure, if for some i, we have  $|M_i| = k + 1 + p$  and  $v \in M_i$ , then v
- these properties. Since we have stopped our procedure, if for some *i*, we have  $|M_i| = k + 1 + p$  and  $v \in M_i$ , then *v* has neighbors in every color class of size at most k + p 1. By (i), the number of color classes of size at least k + p. (including  $M_i$ ) is at most (n - kA - 1)/n = (k + r - 1)/n (here -1 arises because  $|M_i| = k + 1 + r$ ). It follows
- 21 (including  $M_i$ ) is at most  $(n k\Delta 1)/p = (k + r 1)/p$  (here -1 arises, because  $|M_i| = k + 1 + p$ ). It follows that v is adjacent to vertices in at least  $\Delta (k + r 1)/p$  color classes.

### 23 **3. Background for the proof of Theorem 1**

- In this section, we do preparatory work for the proof of Theorem 1: introduce some notions and prove Theorem 1 for  $n \leq 8.8\Delta$ .
- Let G = (V, E) be a graph with maximum degree  $\Delta$  and f be a vertex coloring of G with color classes  $M_1, \ldots, M_{\Delta}$ (some color classes can be empty). For a set  $Y_0$  of color classes of f and a subset V' of V, we define the (V', f)-expansion of  $Y_0$  in G as follows.
- 29 Say that a vertex  $w \in V' \bigcup_{M \in Y_0} M$  is a  $Y_0$ -candidate if w has no neighbors in some color class  $M(w) \in Y_0$ . Let  $Y_1$  be the set of color classes of f containing a  $Y_0$ -candidate. Similarly, for  $h \ge 1$ , a vertex  $w \in V' \bigcup_{M \in Y_0 \cup \dots \cup Y_h} M$  is
- 31 a  $Y_h$ -candidate if w has no neighbors in some color class  $M(w) \in Y_h$ . Let  $Y_{h+1}$  be the set of color classes containing a  $Y_h$ -candidate. Finally, the set  $Y = \bigcup_{i=0}^{\infty} Y_i$  will be called the (V', f)-expansion of  $Y_0$  in G.
- 33 By the construction, we have the following.

Claim 3.1. If Y is the (V', f)-expansion of  $Y_0$  in G, then every vertex  $u \in V' - \bigcup_{M \in Y} M$  has a neighbor in every 35  $M \in Y$ .

For each  $M \in Y - Y_0$ , we define an (M, Y)-recoloring as follows. Suppose that  $M \in Y_{h+1}$  for some  $h \ge 0$ . By the definition of  $Y_{h+1}$ , M contains a  $Y_h$ -candidate  $x_{h+1} \in V'$ . Furthermore, for j = h, h - 1, ..., 1, the color class  $M_j = M(x_{j+1})$  contains a  $Y_{j-1}$ -candidate  $x_j \in V'$ . Then an (M, Y)-recoloring of f is the coloring f' that differs from f only at  $x_1, ..., x_{h+1}$ : for every  $x_j, j = h + 1, h, ..., 1$ , we let  $f'(x_j) = M_{j-1} = M(x_j)$ . For different choices of

- $x_1, \ldots, x_{h+1}$ , we get different recolorings, but the following claim holds by the definition.
- 41 **Claim 3.2.** If  $h \ge 0$  and  $M \in Y_{h+1}$  for some (V', f)-expansion Y of  $Y_0$  in G and f' is an (M, Y)-recoloring of f, then (a) the sizes of almost all color classes in f and f' are the same, only the size of M decreases by one and the size of 43 one color class in  $Y_0$  increases by one;
  - (b) the colors of vertices in V(G) V' are the same in f and f'.  $\Box$

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**1** Lemma 1. Let  $\Delta \ge 3$  and G be a  $K_{\Delta+1}$ -free n-vertex graph with maximum degree  $\Delta$  and

 $|E(G)| \leq n\Delta/10.$ 

3 If  $n/\Delta \leq 8.8$ , then G has an equitable  $\Delta$ -coloring.

**Proof.** Let *G* be an inclusion minimal counterexample to the lemma, and *e* be an edge adjacent to a vertex *v* of the lowest positive degree  $\alpha$ . If  $\alpha > 0.4\Delta$ , then by (1) and the choice of *v*, the cardinality of the set *V'* of non-isolated vertices in *G* is at most  $2|E(G)|/\alpha < 2 \cdot 0.1n\Delta/0.4\Delta = n/2$ . Applying Lemma 2 to G[V'], we obtain a  $\Delta$ -coloring of

- 7 G[V'] in which every color class has at most  $\lceil n/\Delta \rceil$  vertices and each but one color class has less than  $\lceil n/\Delta \rceil$  vertices. Now we can add the remaining isolated vertices to these color classes in order to get the size of every class equal to
- 9  $\lceil n/\Delta \rceil$  or  $\lfloor n/\Delta \rfloor$ . This contradicts the choice of *G*. Therefore,

$$\deg(v) \leq 0.4\Delta$$
.

(2)

- By the minimality of G, G-e has an equitable  $\Delta$ -coloring f. We may assume that the color classes of f are  $M_1, \ldots, M_{\Delta}$ and that  $v \in M_{\Delta}$ . By the definition, the size of every  $M_i$  is either t or t - 1, where  $t = \lfloor n/\Delta \rfloor$ .
- 13 Let  $M_0 = M_{\Delta} v$ . Then  $f_0 = (M_0, \dots, M_{\Delta-1})$  is a proper coloring of  $G_0 = G v$ . Let  $Y_0$  denote the set of color classes in  $f_0$  of size  $|M_0|$ . If some  $M_j \in Y_0$  contains no neighbors of v, then  $|Y_0| > 1$  and hence  $|M_0| = t 1$ . In
- this case, we color  $v_i$  with  $M_j$  and get an equitable  $\Delta$ -coloring of G, a contradiction. Thus, every  $M_j \in Y_0$  contains a neighbor of v. Let  $Y = \bigcup_{j=0}^{\infty} Y_j$  be the  $(V(G_0), f_0)$ -expansion of  $Y_0$  in  $G_0$  (defined at the beginning of the section)
- 17 and y = |Y|. Suppose that for some  $h \ge 0$ , a color class  $M_{h+1}$  in  $Y_{h+1}$  does not contain a neighbor of v. Let f' be an (M, Y)-

19 recoloring of  $f_0$ . By Claim 3.2, if we additionally color v with  $M_{h+1}$ , then we obtain an equitable  $\Delta$ -coloring of G, a contradiction. Thus, every color class in Y contains a neighbor of v and therefore  $y \leq \deg_G(v)$ .

- 21 Let  $V^+ = V(G) \bigcup_{M \in Y} M$ .
- Case 1:  $|M_0| = t 1$ . By the definition of  $Y_0$ , every color class outside of Y has size t. Therefore,  $|V^+| = 23 \quad t(\Delta y) \ge n(\Delta y)/\Delta$ . By (1) and Claim 3.1, we have

$$\frac{n\Delta}{10} \ge |E(G)| \ge y(\Delta - y)\frac{n}{\Delta}.$$
(3)

- 25 For  $\lambda = y/\Delta$ , (3) gives  $\lambda^2 \lambda + 0.1 \ge 0$ . It follows that either  $\lambda > 0.88$  or  $\lambda < 0.12$ . The former contradicts (2), so  $y < 0.12\Delta$ . Since every vertex in  $V^+$  is adjacent to  $M_0$ , we get  $\Delta(t-1) \ge t(\Delta y)$ , i.e.,  $ty \ge \Delta$ . This yields  $t > \frac{1}{0.12} > 8$ ,
- 27 which means  $n > 8\Delta$  and  $t \ge 9$ . Furthermore, since every color class outside of *Y* has size *t* and  $y < 0.12\Delta$ , we get

$$n \ge t(\varDelta - y) + (t - 1)y \ge 9\varDelta - y \ge 8.88\varDelta.$$

- 29 Case 2:  $|M_0| = t 2$ . Recall that no other color class has size less than t 1. Suppose that some  $M \in Y$ , say,  $M \in Y_h$  has t vertices. Then any (M, Y)-recoloring f' of  $f_0$  satisfies the conditions of Case 1 with  $M_h x_h$  in place of  $M_0$ . Since
- 31 Case 1 is proved, we can assume that every  $M \in Y$  has at most t 1 vertices. It follows that the average size of color classes outside of Y is higher than in Y, and therefore higher than  $n/\Delta$ . Thus, (3) holds, and we get  $y < 0.12\Delta$  exactly as
- in Case 1. Similarly to Case 1, we obtain  $\Delta(t-2) \ge (t-1)(\Delta y)$ , i.e.,  $(t-1)y \ge \Delta$ . It follows that  $t-1 > \frac{1}{0.12} > 8$ . Since t-1 is an integer, we have  $n/\Delta > t-1 \ge 9$ .  $\Box$

#### 35 4. The bounded-size version of Theorem 1

An *l-bounded* coloring of a graph *G* is a proper vertex coloring of *G* in which the size of each color class is at most *l*. Clearly, every equitable *k*-coloring of an *n*-vertex graph is [n/k]-bounded, but not every [n/k]-bounded *k*-coloring of an *n*-vertex graph is equitable. Thus, the theorem below is a weaker statement than Theorem 1.

39 **Theorem 3.** Let  $\Delta$ ,  $n \ge 46$ . Suppose that an n-vertex graph G = (V, E) has maximum degree at most  $\Delta$  and  $|E| \le \Delta n/5$ . If G does not contain  $K_{\Delta+1}$ , then G has an  $\lceil n/\Delta \rceil$ -bounded coloring with  $\Delta$  colors.

5

(1)

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- 1 **Proof.** Let *G* be a  $K_{\Delta+1}$ -free graph with maximum degree  $\Delta \ge 46$  and average degree at most  $\Delta/5$ . Let  $t = \lceil n/\Delta \rceil$ . Let
  - $V_l = \{v_1, \ldots, v_l\}$  be the set of vertices of degree at least  $4\Delta/5$ . For  $i = l + 1, \ldots, n$ , consider the following procedure:
- 3 (1) let  $\Delta_i$  be the maximum degree in  $G V_{i-1}$ ;
- (2) if  $\Delta_i \ge 2\Delta/5$ , then let  $v_i$  be a vertex of degree  $\Delta_i$  in  $G V_{i-1}$ ;
- 5 (3) if  $\Delta_i < 2\Delta/5$ , then let  $v_i$  be a vertex of maximum degree in G among vertices in  $G V_{i-1}$ ;
- (4) let  $V_i = V_{i-1} \cup \{v_i\}$  and  $G_i = G[V_i]$ .
- 7 Note that, in general, the resulting ordering is not unique.

**Claim 4.1.** Let *m* be the largest index *i* such that  $\Delta_i \ge 2\Delta/5$ . Then  $m \le \lfloor n/4 \rfloor$ .

9 **Proof.** By the definition, every  $v_i$  for  $l + 1 \le i \le m$  has at least  $2\Delta/5$  adjacent vertices  $v_j$  with j > i. Therefore, *G* contains at least  $(m - l)2\Delta/5$  edges not incident with  $v_1, \ldots, v_l$ . Thus, if  $\Delta/4 < m$ , then

11 
$$|E(G)| \ge l \cdot \frac{4\Delta}{5 \cdot 2} + (m-l) \cdot \frac{2\Delta}{5} > \frac{n}{4} \frac{2\Delta}{5} = \frac{\Delta n}{10},$$

a contradiction.  $\Box$ 

- 13 Suppose that  $m = k(\Delta + 1) + r$  with  $0 \le r \le \Delta$ . Then  $G_m$  has a proper  $\Delta$ -coloring  $f_m = (M_1^m, \dots, M_{\Delta}^m)$  satisfying Corollary 2. Note that by Claim 4.1,  $k + \lceil m/\Delta + 1 \rceil = \lfloor m/\Delta + 1 \rfloor + \lceil m/\Delta + 1 \rceil < t$ . Therefore,  $f_m$  is a *t*-bounded 15 coloring of  $G_m$ .
- We will now complete  $f_m$  to a *t*-bounded  $\Delta$ -coloring of G by constructing consecutively colorings  $f_i$  of  $G_i$  for i = 1 + m, 2 + m, ..., n, in such a way that

$$f_i(v) = f_m(v)$$
 for every  $v \in V_m$ . (4)

- 19 Observe that  $f_m$  satisfies (4). Now, suppose that  $m + 1 \le i \le n$  and  $G_{i-1}$  has a *t*-bounded coloring  $f_{i-1}$  satisfying (4). We will construct  $f_i$  for  $G_i$ .
- 21 Let  $M_1, \ldots, M_{\Delta}$  be the color classes of  $f_{i-1}$ . Let  $Y_0$  denote the set of color classes of cardinality less than *t*. If some  $M_i \in Y_0$  contains no neighbors of  $v_i$ , then we color  $v_i$  with  $M_i$  and have a *t*-bounded coloring  $f_i$  satisfying (4).
- 23 Otherwise, let  $Y = \bigcup_{j=0}^{\infty} Y_j$  be the  $(V_{i-1} V_m, f_{i-1})$ -expansion of  $Y_0$  in  $G_{i-1}$  (defined in Section 3) and y = |Y|. If for some  $h \ge 0$ , a color class  $M_{h+1}$  in  $Y_{h+1}$  does not contain a neighbor of v, then consider an (M, Y)-recoloring
- 25 f' of  $f_{i-1}$ . By Claim 3.2(a), if we additionally color  $v_i$  with  $M_{h+1}$ , then we obtain a *t*-bounded  $\Delta$ -coloring of  $G_i$ . Moreover, by Claim 3.2(b), this new coloring also satisfies (4), as required. Thus, we may assume that every color class
- in Y contains a neighbor of v. This together with the definition of  $Y_0$  and Claim 3.1 yields the following.

Claim 4.2. If  $G_i$  has no t-bounded coloring  $f_i$  satisfying (4), then the  $(V_{i-1} - V_m, f_{i-1})$ -expansion Y of  $Y_0$  in  $G_{i-1}$  possesses the following properties:

- (a) every color class in Y contains a neighbor of  $v_i$  and thus  $y \leq \deg_G(v_i)$ ,
- 31 (b) every vertex  $u \in V_{i-1} V_m \bigcup_{M \in Y} M$  has a neighbor in every  $M \in Y$ ,
  - (c) every color class outside of Y has t vertices.
- 33 Let

$$V^- = V_m - \bigcup_{M \in Y} M$$

35 and

$$V^+ = V_{i-1} - V_m - \bigcup_{M \in Y} M = V_{i-1} - \bigcup_{M \in Y} M - V^-.$$

37 **Claim 4.3.** Using the notation above,  $|V^-| \leq 3n/8\Delta(\Delta - y)$ .

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1 **Proof.** Recall that  $m = k(\varDelta + 1) + r$  with  $0 \le r \le \varDelta$ . By Corollary 2(ii),

$$V^{-}| \leqslant k + \left\lceil \frac{m}{\varDelta + 1} \right\rceil (\varDelta - y) \leqslant \frac{n}{4\varDelta} + \left\lceil \frac{n}{4\varDelta} \right\rceil (\varDelta - y).$$

3 Since  $\Delta \ge 46$  and i > l, we have  $\Delta - y \ge \Delta/5 > 9$ . By Lemma 1,  $n/4\Delta \ge 2.2$  and therefore,  $\lceil n/4\Delta \rceil \le \frac{3}{2.2} \cdot n/4\Delta$ . It follows that

$$_{5} \qquad |V^{-}| \leqslant (\varDelta - y) \left( \frac{n}{4\varDelta(\varDelta - y)} + \frac{3}{2.2} \cdot \frac{n}{4\varDelta} \right) \leqslant \frac{n}{4\varDelta} \left( \varDelta - y \right) \left( \frac{1}{10} + \frac{3}{2.2} \right) < \frac{3n}{8\varDelta} \left( \varDelta - y \right).$$

This proves the claim.  $\Box$ 

7 **Claim 4.4.** The size y of Y is less than  $0.15\Delta$ .

**Proof.** By Claim 4.2(b), at least  $y|V^+|$  edges connect  $V^+$  with  $\bigcup_{M \in Y} M$ . Recall that every  $v_q$  for  $l + 1 \le q \le m$  has at least  $2\Delta/5$  adjacent vertices  $v_j$  with j > q. Thus, at least  $0.4\Delta|V^-|$  edges of *G* are incident with  $V^-$  and hence

$$|E(G)| \ge y|V^+| + \frac{y}{2}|V^-|.$$

11 Since  $|V(G) - \bigcup_{M \in Y} M| \ge t(\Delta - y) \ge n/\Delta(\Delta - y)$ , we have

$$|E(G)| \ge y\left((\varDelta - y)\frac{n}{\varDelta} - |V^-|\right) + \frac{y}{2}|V^-| \ge y\left((\varDelta - y)\frac{n}{\varDelta} - \frac{y}{2}|V^-|\right).$$

13 This and Claim 4.3 yield

$$\frac{n\Delta}{10} \ge y \left(\frac{n}{\Delta} \left(\Delta - y\right) - \left(\Delta - y\right) \frac{3n}{16\Delta}\right) = \frac{13}{16} yn \left(1 - \frac{y}{\Delta}\right).$$

- 15 Denoting  $\lambda = y/\Delta$  and dividing both parts by  $n\Delta$ , we obtain  $\frac{1}{10} \ge \frac{13}{16}\lambda(1-\lambda)$ . Solving this inequality, we get  $\lambda > 0.85$  or  $\lambda < 0.15$ . Since  $i > m \ge l$ , we have  $y \le \deg(v_i) < 0.8\Delta$ . We conclude that  $\lambda = y/\Delta < 0.15$ .  $\Box$
- 17 **Claim 4.5.** The size of  $V^+$  is greater than  $\frac{2n}{3}$ .

**Proof.** Assume that  $|V^+| \leq \frac{2n}{3}$ . As in the proof of Claim 4.4, at least  $y|V^+|$  edges connect  $V^+$  with  $\bigcup_{M \in Y} M$  and at least  $0.4\Delta |V^-|$  edges are incident with  $V^-$ . Hence

$$|E(G)| \ge y|V^+| + 0.4\Delta \left( (\Delta - y) \frac{n}{\Delta} - |V^+| \right) \ge |V^+|(y - 0.4\Delta) + 0.4\frac{n}{\Delta}\Delta(\Delta - y).$$

21 Recall that  $y < 0.15 \Delta$  by Claim 4.4 and  $|V^+| \leq \frac{2n}{3}$ . Thus the last inequality yields

$$\frac{n\varDelta}{10} > \frac{2n}{3} \left( y - 0.4\varDelta \right) + 0.4n(\varDelta - y).$$

23 Dividing both parts by  $n\Delta$ , we obtain

$$\frac{1}{10} > \frac{2}{3}\frac{y}{4} - \frac{8}{30} + \frac{2}{5} - \frac{2y}{54}$$

25 which is false for  $y/\Delta \ge 0$ . This contradiction proves the claim.  $\Box$ 

**Claim 4.6.** Let  $M_1$  be a color class of the smallest size in Y. Then

$$|M_1| \leqslant \frac{n+y}{\varDelta} - 1. \tag{5}$$

**Proof.** Since  $v_i$  is not colored,  $|M_1| < n/\Delta$ . If  $|M_1| \le t - 2$ , the conclusion is obvious. Suppose that  $|M_1| = t - 1$ . 29 Since every color class not in *Y* has size *t*, by the minimality of  $|M_1|$ ,  $n + y \ge t\Delta$ . This proves the claim.  $\Box$ 

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1 Now we are ready to finish the proof. Define  $z = 4n/9\Delta + \frac{17}{30}$ . *Case* 1:  $|M_1 \cap V_m| \le z$ . Then the number of neighbors of  $M_1$  is at most  $z\Delta + (|M_1| - z)(2\Delta/5) = z(3\Delta/5) + |M_1|2\Delta/5$ .

3 By Claims 4.6 and 4.4, this is less than

$$\left(\frac{4n}{9\Delta} + \frac{17}{30}\right)\frac{3\Delta}{5} + \left(\frac{n}{\Delta} - 0.85\right) \cdot \frac{2\Delta}{5} = \frac{2n}{3}.$$
(6)

5 Since every vertex in  $V^+$  is a neighbor of  $M_1$ , (6) contradicts Claim 4.5.

Case 2:  $|M_1 \cap V_{\mu}| = z + x$ , where x > 0. We will prove that  $M_1$  has at most  $z(3\Delta/5) + |M_1|(2\Delta/5)$  neighbors in 7  $V^+$ , which by (6) would give the same contradiction as in Case 1.

By Corollary 2(iii), every vertex of  $M_1 \cap V_m$  has at most (k + r - 1)/(z + x - k - 1) neighbors in  $V^+$ . Thus, it is 9 enough to prove that

$$\frac{k+r-1}{z+x-k-1}(z+x) \leqslant z\varDelta + x\,\frac{2\varDelta}{5},\tag{7}$$

since the RHS of (7) is the maximum amount contributed by z + x vertices to  $z\Delta + (|M_1| - z)(2\Delta/5)$  in Case 1. Note that (7) is equivalent to

13 
$$\left(z\varDelta + x\frac{2\varDelta}{5}\right)(z+x-k-1) - (k+r-1)(z+x) \ge 0.$$
 (8)

For x = 0, (8) becomes

15 
$$z \Delta (z+k-1) - z(k+r-1) \ge 0$$
,

which reduces to

21

29

$$17 z \ge \frac{k+r-1}{\varDelta} + k + 1. (9)$$

Recall that  $m = k(\varDelta + 1) + r \leq \lfloor n/4 \rfloor$ , i.e.,  $n \geq 4(k\varDelta + k + r)$ , and that  $z = 4n/9\varDelta + \frac{17}{30}$ . It follows that

19 
$$z \ge \frac{16k}{9} + \frac{16}{9}\frac{k+r}{4} + \frac{17}{30},$$
 (10)

which yields (9) for  $k \ge 1$ . This proves (8) for x = 0.

Now consider the LHS of (8) as a function g(x). Then

$$g'(x) = \frac{2\Delta}{5}(z+x-k-1) + \left(z\Delta + x\frac{2\Delta}{5}\right) - (k+r-1) = \frac{4\Delta}{5}x + \frac{7\Delta}{5}z - \frac{2\Delta}{5}(k+1) - (k+r-1).$$

23 We want to show that  $g'(x) \ge 0$  for every x > 0. This would prove (8) and thus the theorem. By the last equality, the condition  $g'(x) \ge 0$  is equivalent to

25 
$$\frac{7}{5}z \ge \frac{2}{5}(k+1) - \frac{k+r-1}{\Delta}$$

This inequality is implied by (10) for  $k \ge 1$ .  $\Box$ 

#### 27 5. Proof of Theorem 1

The algorithm in the previous section produces a *t*-bounded  $\Delta$ -coloring of *G*, but this coloring might have 'small' color classes. In order to 'correct' the coloring, we use a slight variation of the technique used above.

- Consider *t*-bounded colorings of G obtained in the course of proof of Theorem 3. In particular, each vertex  $v \in V(G) V_m$  has at most 0.8 $\Delta$  neighbors in G and at most 0.4 $\Delta$  neighbors in  $G V_m$ . Among such colorings with a
- fixed coloring  $f_m$  of  $G[V_m]$  satisfying Corollary 2, choose a coloring  $f_0$  with fewest color classes of size t. We will prove that  $f_0$  has no color classes of size t = 2 or less
- 33 prove that  $f_0$  has no color classes of size t 2 or less.

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(11)

- 1 Let  $Y_0$  be the set of color classes of size at most t 2 and assume that  $Y_0$  is non-empty. Let  $Y = \bigcup_{j=0}^{\infty} Y_j$  be the  $(V(G) V_m, f_0)$ -expansion of  $Y_0$  in G and y = |Y|.
- 3 **Claim 5.1.** *Y possesses the following properties:*
- (a) every vertex  $u \in V(G) V_m \bigcup_{M \in Y} M$  has a neighbor in every  $M \in Y$ ,
- 5 (b) every color class in Y has at most t 1 vertices,
  - (c) every color class outside of Y has at least t 1 vertices.
- 7 **Proof.** Claim 3.1 implies (a), and the definition of  $Y_0$  yields (c). To prove (b), assume by contradiction that for some  $h \ge 0$ , a color class  $M_{h+1}$  in  $Y_{h+1}$  has cardinality t. Consider an (M, Y)-recoloring f' of  $f_0$ . By Claim 3.2(a), f'
- 9 is a *t*-bounded  $\Delta$ -coloring of G with fewer color classes of size *t*. Moreover, by Claim 3.2(b), f' satisfies (4). This contradicts the choice of  $f_0$ .  $\Box$

11 Since there is a color class M' of size  $t, y < \Delta$ . Since every vertex in  $M' - V_m$  has neighbors in each color class of Y,

$$y \leq 0.8 \Delta$$
.

13 Let

$$V^- = V_m - \bigcup_{M \in Y} M$$

15 and

$$V^+ = V - V_m - \bigcup_{M \in Y} M = V - \bigcup_{M \in Y} M - V^-$$

17 Now, Claim 4.3 holds: the proof simply repeats that in the previous section. By Claim 5.1(b) and (c),

$$\left| V(G) - \bigcup_{M \in Y} M \right| \ge (\Delta - y) \frac{n}{\Delta}.$$
(12)

19 Thus we can essentially repeat the proofs of Claims 4.4 and 4.5, and conclude that they hold for our *Y*. Let  $M_1 \in Y_0$ . By the definition of  $Y_0$ ,  $|M_1| \le t - 2 < n/\Delta - 1$ , which is stronger than Claim 4.6. Therefore, all the

21 calculations in the previous section following Claim 4.6 go through and we get a contradiction to our assumption.

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