

Packing and Covering Triangles in K_4 -free Planar Graphs*

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Abstract

We show that every K_4 -free planar graph with at most ν edge-disjoint triangles contains a set of at most $\frac{3}{2}\nu$ edges whose removal makes the graph triangle-free. Moreover, equality is attained only when G is the edge-disjoint union of 5-wheels plus possibly some edges that are not in triangles. We also show that the same statement is true if instead of planar graphs we consider the class of graphs in which each edge belongs to at most two triangles. In contrast, it is known that for any $c < 2$ there are K_4 -free graphs with at most ν edge-disjoint triangles that need more than $c\nu$ edges to cover all triangles.

1 Introduction

The main motivation for this paper is an old conjecture of Tuza about packing and covering of triangles by edges. A *triangle packing* in a graph G is a set of pairwise edge-disjoint triangles. A *triangle edge cover* in G is a set of edges meeting all triangles. We denote by $\nu(G)$ the maximum cardinality of a triangle packing in G , and by $\tau(G)$ the minimum cardinality of a triangle edge cover for G . It is clear that for every graph G we have $\nu(G) \leq \tau(G) \leq 3\nu(G)$.

In 1984, Tuza [11] proposed the following conjecture.

Conjecture 1 *For every graph G , $\tau(G) \leq 2\nu(G)$.*

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The complete graphs K_4 and K_5 show that this bound is tight. The conjecture is known to be true for certain special classes of graphs, for example K_5 -free chordal graphs and planar graphs (Tuza [12]), more generally, graphs without a subdivision of $K_{3,3}$ (Krivelevich [8]), tripartite graphs (Haxell and Kohayakawa [6]), odd-wheel-free graphs and four-colourable graphs (Aparna Lakshmanan, Bujtás and Tuza [1]). Weighted versions of the problem were studied in [2]. The only general bound known [5] shows that $\tau(G) \leq \frac{66}{23}\nu(G)$ for every graph G .

Our aim in this paper is to study the planar case more closely. As just mentioned, Tuza [12] proved that the conjecture is true for planar graphs, and it is tight for K_4 . In [3] it was shown that equality holds if and only if G is an edge-disjoint union of copies of K_4 plus possibly some edges that are not in triangles. Here we consider the next step.

Theorem 1 *Let G be a K_4 -free planar graph. Then $\tau(G) \leq \frac{3}{2}\nu(G)$, and equality holds if and only if G is an edge-disjoint union of 5-wheels plus possibly some edges that are not in triangles.*

Our main tool will be a similar result for a different class of graphs. A graph G is *flat* if each edge of G belongs to at most two triangles. Observe that a planar graph is flat if it does not have separating triangles. Flat graphs can be far from planar, but the result we prove for them is the same:

Theorem 2 *Let G be a K_4 -free flat graph. Then $\tau(G) \leq \frac{3}{2}\nu(G)$, and equality holds if and only if G is an edge-disjoint union of 5-wheels plus possibly some edges that are not in triangles.*

It is worth mentioning that excluding K_4 does not have a similar effect in the general case of Tuza's Conjecture: for every $\epsilon > 0$ there exists a K_4 -free graph G_ϵ satisfying $\tau(G_\epsilon) > (2 - \epsilon)\nu(G_\epsilon)$ (see [7]).

Our proof of Theorem 1 makes use of some special properties of the *triangle graph* $T = T(G)$, defined as follows: the vertices of T are the triangles of G , and two vertices are adjacent if the corresponding triangles have an edge in common. These properties are established in Section 2 for flat graphs. Then in Section 3 we prove Theorem 1.

2 Triangle graphs of flat graphs

For every flat graph G , each edge of the triangle graph $T = T(G)$ naturally corresponds to an edge of G . Moreover $T(G)$ is *subcubic* (i.e. has maximum degree at most three), the parameter $\nu(G)$ is equal to the independence number $\alpha(T)$ of T , and $\tau(G)$ is the minimum size $\beta'(T)$ of an edge cover of the vertices of T . If T has no isolated vertices then by Gallai's Theorem, $\beta'(T) + \alpha'(T) = |V(T)|$, where $\alpha'(T)$ denotes the maximum size of a matching in T . Thus to get information on packing and covering triangles in G , we will start by studying $\alpha(T)$ and $\alpha'(T)$.

Let \mathcal{F} denote the family of all triangle graphs for flat K_4 -free graphs. Graphs in \mathcal{F} have some interesting properties.

Claim 3 *If triangles B_1, B_2, B_3 of a flat graph G form a triangle in the triangle graph $T(G)$, then $|V(B_1 \cup B_2 \cup B_3)| = 4$.*

Proof. Since B_1 and B_2 share an edge, $|V(B_1 \cup B_2)| = 4$. If B_3 has a vertex not in $B_1 \cup B_2$, then it shares the same edge with both B_1 and B_2 . But then this edge witnesses that G is not flat. \square

Claim 4 *Each triangle graph $T \in \mathcal{F}$ has no K_3 -subgraphs.*

Proof. Suppose that triangles B_1, B_2, B_3 induce a K_3 in T . Then by Claim 3, all these triangles are contained in the same 4-element set, say M . But then all the 6 pairs of vertices of M must be edges, and hence $G[M]$ is a K_4 , a contradiction to the definition of \mathcal{F} . \square

Claim 5 *Let G be a flat K_4 -free graph. Then any 5-cycle (B_1, \dots, B_5) in $T(G)$ corresponds to a 5-wheel in G .*

Proof. Since B_2 shares two vertices with B_1 and two vertices with B_3 , there is a vertex v_0 that belongs to all three triangles. So we have the situation in Fig. 1. Triangle B_4 must share with B_3 either v_3v_4 or v_0v_4 . If it shares v_3v_4 , then since G is K_4 -free, its third vertex, say v_5 , is not v_1 . Then B_5 needs to share two vertices with each of the two disjoint triangles, B_1 and B_4 , an impossibility. So, $\{v_0, v_4\} \subset B_4$. Since $B_1B_4 \notin E(T(G))$, $v_5 \neq v_1$. Then the only way that a triangle B_5 shares an edge with B_1 and an edge with B_4 is that $v_0 \in B_5$. Since $B_5B_2, B_5B_3 \notin E(T(G))$, the other two vertices of B_5 are v_1 and v_5 . \square

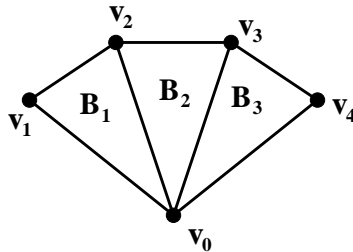


Figure 1:

To study the independence number of graphs in \mathcal{F} we make use of the following theorem of Fajtlowicz [4] and Stanton [10].

Theorem 6 *Every triangle-free subcubic graph with n vertices has an independent set of size at least $5n/14$.*

Remark 1. Let Q denote the graph shown in Fig. 2. It is routine to check that Q is the unique triangle-free subcubic graph with 11 vertices and independence number 4 (see Appendix).

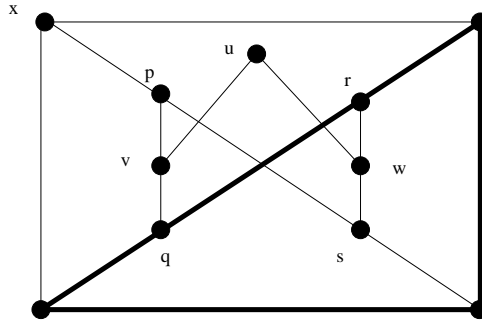


Figure 2: Graph Q .

Lemma 7 *Let n be odd and H be a triangle-free subcubic graph with n vertices. Then $\alpha(H) \geq (n+1)/3$. Moreover if equality holds then H is a 5-cycle or a copy of Q .*

Proof. We note that $5n/14 > (n+1)/3$ for $n \geq 15$, so by Theorem 6 we only need to consider odd $n \leq 13$. It is easy to check that the inequality $\lceil 5n/14 \rceil \geq (n+1)/3$ also holds for these values (for example, $\lceil \frac{5 \cdot 13}{14} \rceil = 5 > \frac{13+1}{3}$), and that equality occurs only for $n \in \{5, 11\}$. The only 5-vertex triangle-free graph with independence number 2 is C_5 . By Remark 1, if $n = 11$ and $\alpha(H) = 4$, then $H = Q$. \square

Lemma 8 *Let H be a triangle-free subcubic graph without isolated vertices that does not contain Q . Then $\beta'(H) \leq \frac{3}{2}\alpha(H)$. Moreover, if equality holds then H is a disjoint union of C_5 's.*

Proof. Since $\beta'(H) = n - \alpha'(H)$ where $n = |V(H)|$, it suffices to prove that $n - \alpha'(H) \leq \frac{3}{2}\alpha(H)$. Now if H has a perfect matching then $\alpha'(H) = n/2$ and the result follows immediately from Theorem 6. Thus we may assume that H has no

perfect matching. Let S be a *Tutte set*, that is, a subset of $V(H)$ such that $H - S$ has $c = |S| + n - 2\alpha'(H)$ odd components. If we let S be a set of maximal size with respect to this property then every component of $H - S$ is odd and hypomatchable. Let H_i , $1 \leq i \leq c$ denote the components of $H - S$, and set $n_i = |V(H_i)|$.

Now $\alpha(H) \geq \sum_{i=1}^c \alpha(H_i) \geq \sum_{i=1}^c (n_i + 1)/3$ (using Lemma 7), and so we obtain

$$\alpha(H) \geq (n - |S| + c)/3 = (2n - 2\alpha'(H))/3.$$

This implies the first assertion of the lemma. To show the second assertion, assume equality holds. Since H contains no copy of Q we know that $\sum_{i=1}^c \alpha(H_i) = \sum_{i=1}^c (n_i + 1)/3$ implies that every component of $H - S$ is a 5-cycle. If $S = \emptyset$, then H is the disjoint union of c 5-cycles, as claimed. Suppose that $x \in S$. Since H is triangle-free, x has at most two neighbors in each of the 5-cycles, H_1, \dots, H_c . So, in each H_i we can choose two non-neighbors y_i, y'_i of x such that $y_i y'_i \notin E(H)$. Since the set $\{x\} \cup \bigcup_{i=1}^c \{y_i, y'_i\}$ is independent, we have $\alpha(H) \geq 1 + 2c > \sum_{i=1}^c (n_i + 1)/3$, a contradiction to our assumption of equality. This completes the proof. \square

Proof of Theorem 2. By Claim 4 we know that $T(G)$ is triangle-free, and we may assume that $T(G)$ has no isolated vertices. We claim that $T(G)$ does not contain Q , in which case the proof is complete by Lemma 8 and Claim 5. Suppose on the contrary that $T(G)$ does contain Q . Observe that Q contains a 5-cycle C (marked in bold in Fig. 2) together with another vertex x that is adjacent to two non-adjacent vertices of C (see Fig. 2). By Claim 5 we know that C corresponds to a 5-wheel W in G . Then since G is flat, it is not possible for another triangle x to share an edge with two non-adjacent triangles in W , as each must be adjacent to x along a rim edge. Thus $T(G)$ cannot contain Q , as required. \square

3 Planar graphs

We are now ready to prove Theorem 1. Let G be a counter-example with the smallest $|V(G)| + |E(G)|$. For graphs with at most 4 vertices the statement is evident, so $|V(G)| \geq 5$. Also by the minimality we have:

- (S1) *Each edge of G belongs to a triangle;*
- (S2) *For each triangle B , at least 2 edges of B belong to other triangles.*

By Theorem 2, G is not flat. This means that there is a separating triangle B_0 such that some other triangles are inside B_0 and some outside. We choose B_0 to be a separating triangle with minimum interior, so that in particular the subgraph of G induced by the vertices inside and on B_0 is a flat plane graph without separating triangles (and B_0 is one of its triangles). We denote this subgraph by G_1 and the

graph $G_1 - E(B_0)$ by G'_1 . Similarly the "outside" subgraph $G - G'_1$ will be denoted by G_2 , and $G - G_1 = G_2 - E(B_0)$ will be denoted by G'_2 .

For $i = 1, 2$, let $\nu_i = \nu(G'_i)$ and $\tau_i = \tau(G'_i)$. We now derive some properties of G .

(S3) *Let W be a 5-wheel in G and (x, y, z) be a 3-face of W . Then it cannot happen that all vertices of G not in W are inside (x, y, z) while all of $W - \{x, y, z\}$ are outside (x, y, z) .*

Proof. Let x be the vertex of degree 5 in W and let the rim of W be (y, u, v, w, z) . Suppose that only the face (x, y, z) contains vertices of $G - W$. Let $G' := G - u - v - w$. Then $\nu(G' - xy) \leq \nu(G) - 2$, since any set of edge-disjoint triangles in $G - xy$ can be complemented by (x, y, u) and (x, v, w) . So, if $T(G' - xy)$ is not a disjoint union of C_5 s, then by the minimality of G , $\tau(G' - xy) < 3(\nu(G) - 2)/2$, and we can add 3 edges including xy (e.g. $\{xy, xv, xw\}$) that cover all remaining triangles of G . Thus, $T(G' - xy)$ is the disjoint union of C_5 s. If one of the corresponding 5-wheels in G contains the edge xz , then we can choose a covering of this wheel with 3 edges using xz and cover the triangles of W not containing xz with two edges, one of which is xy . So, xz does not belong to triangles in $G' - xy$. By symmetry, we find that $T(G' - xz)$ is the disjoint union of C_5 s and xy does not belong to triangles in $G' - xz$. Therefore $T(G' - xy) = T(G' - xz) = T(G' - xy - xz)$ is a disjoint union of C_5 s. Now a similar argument shows that yz does not belong to triangles in $G' - xy - xz$. It follows that $T = T(G)$ is a disjoint union of C_5 s, a contradiction. \square

Everywhere below we assume that $B_0 = (x, y, z)$.

(S4) $\nu(G) = \nu_1 + \nu_2 + 1$.

Proof. Indeed, since the edges in B_0 do not participate in triangles of G'_1 and G'_2 , $\nu(G) \geq \nu_1 + \nu_2 + 1$. On the other hand, suppose $\nu(G) \geq \nu_1 + \nu_2 + 2$. By definition, $\tau(G) \leq \tau_1 + \tau_2 + |E(B_0)| = \tau_1 + \tau_2 + 3$. By the minimality of G , $\tau_1 + \tau_2 \leq \frac{3}{2}(\nu_1 + \nu_2)$ with equality only if all components of $T(G - E(B_0))$ are C_5 . So, we may assume that all components of $T(G - E(B_0))$ are C_5 and that $\nu(G) = \nu_1 + \nu_2 + 2$, since otherwise we are done. By (S1) and (S2), $\nu_1 \geq 1$ and $\nu_2 \geq 1$, and hence each of G'_1 and G'_2 contains at least one 5-wheel. Moreover, each such 5-wheel shares an edge with a triangle containing an edge of B_0 , since otherwise, we can delete the edges of a 5-wheel W and use induction for $G - E(W)$.

Case 1: Every 5-wheel in G'_1 shares at most one edge with triangles containing edges of B_0 . Then we can cover all triangles in G_1 apart from B_0 with $3k$ edges, where k is the number of (edge disjoint) 5-wheels in G'_1 , which is equal to $\nu_1/2$. On the other hand, $\nu(G_2) \leq \nu - \nu_1$ and by the minimality of G , $\tau(G_2) < 1.5(\nu - \nu_1)$, unless $T(G_2)$ is a union of disjoint C_5 . But G_2 is obtained from G'_2 by adding the edges of B_0 , and the edges of a triangle in a 5-wheel (or any two edges) do not cover all triangles in this wheel. So if $T(G'_2)$ is the union of disjoint C_5 , then $T(G_2)$ is not. This finishes the case.

Case 2: Some 5-wheel W in G'_1 shares more than one edge with triangles containing edges of B_0 . By the choice of B_0 , the outside face of W is its 5-face. If a triangle

D in G shares an edge with W and shares an edge with B_0 , then W shares a vertex with B_0 . Furthermore, W cannot share more than one vertex with $B_0 = (x, y, z)$, since otherwise $G[W]$ contains an edge not in W and adding any edge to W creates a K_4 . We may assume that x belongs to W and y and z not. Then there is no triangle containing yz sharing an edge with W . Let the rim of W be (x, x_1, x_2, x_3, x_4) and the center be x_0 . Since the center of W is inside its rim (we call such a 5-wheel *normal*) and G_1 is flat, we may assume that the triangles containing an edge of W and an edge of B_0 are (x, y, x_1) and (x, z, x_4) . The edges yx_1 and zx_4 cannot belong to the same 5-wheel in G'_1 , since it would be normal and yx_1 and zx_4 would be rim edges, so together with yz we would find a K_4 in G . Hence, there is a set F of $\nu_1 + 2$ edge-disjoint triangles in $G_1 - zy$ consisting of triangles (x, y, x_1) , (x, z, x_4) , (x_0, x_1, x_2) , (x_0, x_3, x_4) , and two triangles from each other 5-wheel in $G_1 - zy$ (see Figure 3). It follows that $\nu(G_2 - xy - xz) \leq \nu_2$. Now $T(G'_2)$ is a disjoint union of C_5 , and y and z do not both belong to the same 5-wheel in G'_2 otherwise they form a K_4 . Thus if yz is in a triangle in $G_2 - xy - xz$ then we could find ν_2 triangles in G'_2 avoiding it, giving a total of $|F| + \nu_2 + 1 = \nu_1 + \nu_2 + 3$ edge-disjoint triangles, a contradiction. Therefore yz does not belong to a triangle in G_2 other than B_0 . There is at most one triangle, say (y, z, u) in G_1 distinct from B_0 that contains zy . If there is no such triangle, or (y, z, u) shares an edge with a 5-wheel in G'_1 , then we find a covering containing $1.5\nu_2$ edges in G'_2 , edges xy , xz and $1.5\nu_1$ edges in G'_1 covering all triangles in 5-wheels of G'_1 and the triangle (y, z, u) , if it exists. See Figure 3, where the bold edges are an example of part of a suitable cover. But if (y, z, u) exists and does not share any edge with a 5-wheel in G'_1 , then $\nu(G_1) \geq \nu_1 + 3$, a contradiction. \square

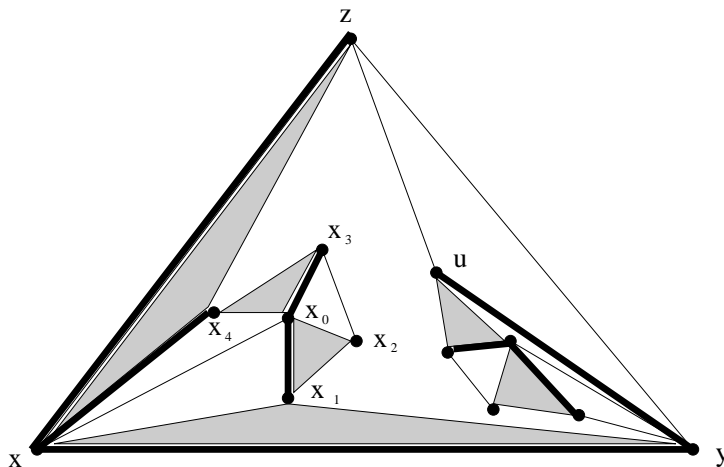


Figure 3: The shaded triangles form an example of the set F .

Property (S4) implies also

(S5) For $i = 1, 2$, $\nu(G_i) = \nu_i + 1$.

(S6) For every $X \subset E(B_0)$, $\nu(G_1 - X) + \nu(G_2 - (E(B_0) - X)) \leq \nu_1 + \nu_2 + 1$.

(S7) For every edge e of B_0 , either $\nu(G'_1 + e) = \nu_1 + 1$ or $\nu(G'_2 + e) = \nu_2 + 1$.

Proof. Assume that $\nu(G'_1 + xy) = \nu_1$ and $\nu(G'_2 + xy) = \nu_2$. If neither of $T(G'_1 + xy)$ and $T(G'_2 + xy)$ is a disjoint union of C_5 s, then we can get a triangle cover for G by adding edges yz and xz to the union of optimum triangle coverings in $G'_1 + xy$ and $G'_2 + xy$. This would imply that

$$\tau(G) \leq 2 + (1.5\nu_1 - 0.5) + (1.5\nu_2 - 0.5) = 1.5(\nu_1 + \nu_2) + 1 = 1.5\nu(G) - 0.5.$$

So, there is $j \in \{1, 2\}$ such that $T(G'_j + xy)$ is a disjoint union of C_5 s.

Case 1: Neither of xz and yz belongs to a triangle in G_j other than B_0 . If xy also does not belong to a triangle in G_j other than B_0 , then each triangle of G belongs either to G'_j or to G_{3-j} which are edge-disjoint. So we are done by the minimality of G . So, xy belongs to a 5-wheel in $G'_j + xy$, and hence there is a set S of $1.5\nu_j$ edges covering all triangles in $G'_j + xy$ such that $xy \in S$. If $T(G_{3-j})$ is not a disjoint union of C_5 s, then there is $S' \subset E(G_{3-j})$ covering all triangles with $|S'| < 1.5(\nu_{3-j} + 1)$; so that $|S \cup S'| < 1.5\nu(G)$. Suppose that $T(G_{3-j})$ is a disjoint union of C_5 s. Since B_0 is a triangle in G_{3-j} , and $xy \in S$, we again need fewer than $1.5\nu_{3-j}$ edges to cover triangles in $G_{3-j} - xy$.

Case 2: The edge xz belongs to a triangle in G_j other than B_0 (say, (x, z, u)), and the edge yz does not. We claim that

$$\nu(G'_j + xz) \geq \nu_j + 1. \quad (1)$$

If (1) does not hold, then some 5-wheel W in $G'_j + xy$ does not have two edge-disjoint triangles that are also disjoint from (x, z, u) and do not contain edge xy . Since $xz \notin E(G'_j + xy)$, W has at most two common vertices with (x, z, u) , otherwise G contains a K_4 . Thus xy is an edge of W and some edge $e \in \{xu, zu\}$ is an edge of W . If $e = zu$, then W contains all of (x, z, u) . So, $e = xu$. Since $W - xy - xu$ has no two edge-disjoint triangles, x is the center of W . Since $G[\{x, y, z, u\}] \neq K_4$, $uy \notin E(G)$. Then there is no face of W to put z so that it is adjacent to all of x, y and u . So, (1) holds.

By (1) and (S6), $\nu(G_{3-j} - xz) = \nu_{3-j}$. In our case (i.e. yz is not in a triangle in G_j other than B_0) we also have $\nu(G_j - xz) = \nu_j$. So we can get a triangle cover for G by adding edge xz to the union of optimum triangle coverings in $G_1 - xz$ and $G_2 - xz$. This yields $\tau(G) \leq 1.5\nu_1 + 1.5\nu_2 + 1 < 1.5\nu(G)$.

Case 3: Edge xz belongs to a triangle (x, z, u) in G_j other than B_0 and the edge yz belongs to a triangle (y, z, v) in G_j other than B_0 . Since G is K_4 -free, $v \neq u$. Similarly to Case 2, we claim that

$$\nu(G_j) \geq \nu_j + 2, \quad (2)$$

which would contradict (S4). If (2) does not hold, then some 5-wheel W in $G'_j + xy$ does not have two edge-disjoint triangles that are also disjoint from (x, z, u) and (y, z, v) . Then W contains an edge $e_1 \in \{xu, uz\}$ and an edge $e_2 \in \{yv, vz\}$, and at least one of them is incident to the center of W . If $e_1 = xu$ then $z \notin W$, otherwise $W + zx$ contains K_4 . Thus $e_2 = yv$ and one of $\{x, u, v, y\}$ is the center of W , but then we find a K_4 containing z . This contradiction shows that $e_1 = uz$, and similarly $e_2 = vz$. Then z is the center of W and $x, y \notin W$. Since $xy \in E(G)$, x and y are in the same face of W . Since the set $\{x, y\}$ is adjacent to z, u , and v , this face should be (u, z, v) . In particular, the triangle (u, z, v) separates (x, z, y) from all other vertices of W . So, by the choice of (x, z, y) , $j = 2$. If a face F of W distinct from (u, z, v) contains a vertex not in W , then by (S1) it contains a triangle distinct from itself. Since $T(G'_j + xy)$ is a disjoint union of C_5 s, F contains a 5-wheel W' distinct from W and hence edge-disjoint from W . So, deleting the edges of W' from G does not destroy other triangles of G , and we can apply induction hypothesis to $G - E(W')$, a contradiction. Thus no face of W apart from (u, z, v) contains vertices. Now we are done by (S3). \square

Final calculations By (S7), there is $j \in \{1, 2\}$ such that $\nu(G'_j + xy) = \nu_j + 1$. Then by (S6), $\nu(G_{3-j} - xy) = \nu_{3-j}$. Then again by (S7), $\nu(G'_j + xz) = \nu_j + 1$ and $\nu(G'_j + zy) = \nu_j + 1$, and again by (S6), $\nu(G_{3-j} - xz) = \nu_{3-j}$ and $\nu(G_{3-j} - zy) = \nu_{3-j}$. By the minimality of G there is an edge-cover S_j of triangles in G_j with $|S_j| \leq 1.5(1 + \nu_j)$ (with equality only if $T(G_j)$ is the disjoint union of C_5 s). Since S_j covers B_0 , we may assume that $xy \in S_j$. By above, there is an edge-cover S_{3-j} of triangles in $G_{3-j} - xy$ with $|S_{3-j}| \leq 1.5\nu_{3-j}$ (with equality only if $T(G_{3-j} - xy)$ is the disjoint union of C_5 s). So, if at least one of $T(G_j)$ and $T(G_{3-j} - xy)$ is not the disjoint union of C_5 s, then we are done. Suppose now that they both are. If the 5-wheels of $G_{3-j} - xy$ contain neither xz nor yz , then $T(G)$ is the disjoint union of C_5 s, and we are done. So, we may assume that xz is in some 5-wheel in $G_{3-j} - xy$. Note that there is an edge-cover S'_j of triangles in G_j with $|S'_j| = 1.5(1 + \nu_j)$ containing xz . Furthermore, if we delete from $G_{3-j} - xy$ the edge xz , we destroy one of the 5-wheels, and when we then add to it the edge xy , the resulting graph $G_{3-j} - xz$ is not the edge-disjoint union of 5-wheels. Since $\nu(G_{3-j} - xz) = \nu_{3-j}$, we are done.

4 Appendix

Here for completeness we show that if G is a triangle-free subcubic graph on 11 vertices with independence number 4, then it is isomorphic to the graph Q shown in Figure 2.

Since G has 11 vertices we know that it has a vertex a with even degree. If $d(a) = 0$ then $G - a$ is a graph with 10 vertices, that contains no triangle and no independent set of size 4, contradicting the fact that the Ramsey number $R(3, 4)$ is 9. Therefore $d(a) = 2$. Let b and c denote the neighbours of a . Then $G' = G - \{a, b, c\}$

has 8 vertices, no triangle and no independent set of size 4. Thus G' is an *extremal Ramsey graph* for $R(3, 4)$, which means that it shows $R(3, 4) \geq 9$. A complete list of such graphs is known (see e.g. [9]). They are: the graph H obtained by removing the vertex u of degree 2 in Q together with its neighbours v and w (see Figure 2), the graph $H_1 = H + pq$, and $H_2 = H_1 + rs$.

Observe that in H_2 (and therefore also in H and H_1), for each $y \in \{p, q, r, s\}$ there exists an independent set S_y of size 3 that contains y and no element of $\{p, q, r, s\} \setminus \{y\}$. Since there must be an edge of G from $\{b, c\}$ to S_y (otherwise G contains an independent set of size 5), and the two vertices in $S_y \setminus \{y\}$ each have degree 3 in H , the only possibility is that $G' = H$ and each of b and c has two neighbours in $\{p, q, r, s\}$. Since ps and qr are edges and G is triangle-free, we conclude that G is isomorphic to Q . \square

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