

LOWER BOUND OF THE HADWIGER NUMBER OF GRAPHS BY THEIR AVERAGE DEGREE

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Received 1 February 1982

The aim of this paper is to show that the minimum Hadwiger number of graphs with average degree k is $O(k/\sqrt{\log k})$. Specially, it follows that Hadwiger's conjecture is true for almost all graphs with n vertices, furthermore if k is large enough then for almost all graphs with n vertices and nk edges.

1. Introduction

Let us recall the notion of the Hadwiger number. The following operations are called elementary contractions:

- substitution of two adjacent points v_1 and v_2 for a new point v_3 connected to the points connected to v_1 or v_2 ;
- removal of an edge;
- removal of a point.

The graph G is said to be contractable to the graph H if H can be obtained from G by means of a sequence of elementary contractions. The Hadwiger number $\eta(G)$ of a graph G is the maximum order of the complete graphs to which G can be contracted.

A contraction of a connected graph $G=(V, E)$ to the complete graph of r vertices also can be considered as an onto mapping $\psi: V \rightarrow \{1, 2, \dots, r\}$ with the following properties:

- every subgraph $\psi^{-1}(i)$ of G is connected ($1 \leq i \leq r$);
- for any integers $1 \leq i < j \leq r$, there exist vertices $v \in \psi^{-1}(i)$ and $w \in \psi^{-1}(j)$ such that $(v, w) \in E$.

The contraction is one of the most natural operations on graphs. Thus, it is very important to study the Hadwiger number and its relation to other features of graphs. For example, let us mention Hadwiger's well-known conjecture ([2]) that the chromatic number $\chi(G)$ of a graph G is not greater than $\eta(G)$. Let $w(k) = \min \{\eta(G) : \chi(G) \cong k\}$. Then Hadwiger's conjecture is equivalent to the following statement: $w(k) = k$ for any natural number k . Wagner [7] showed that $w(k) \cong 4 + \log_2(k - 1/3)$. Mader's [5] result implies that $w(k) > (k/16) \log_2 k$.

Many authors (see e.g. [5], [6], [8], [9]) are interested in how small Hadwiger number a graph $G=(V, E)$ can have if $|E|/|V| \cong k$. Denote this minimum by $\eta(k)$. Considering contraction as mapping, we can easily prove that the Hadwiger number of almost all graphs of n vertices is at most $n/\sqrt{\log n}$. It means that there exists a graph $G=(V, E)$ with $|V|=n$, $|E| \cong n^2/4$ and $\eta(G) \cong n/\sqrt{\log n}$. Thus $\eta(k) \cong 4k/\sqrt{\log k}$ if k is large enough. On the other hand, Mader [5] proved that $\eta(k) > k/8 \log_2 k$. That is, if k is large then

$$\frac{k}{8 \log_2 k} < \eta(k) \cong \frac{4k}{\sqrt{\log k}}.$$

The aim of this paper is to prove that the order of $\eta(k)$ is $O(k/\sqrt{\log k})$. The consequences of this statement improve the lower bound of $w(k)$ and prove Hadwiger’s conjecture for almost all graphs of n vertices and kn edges if k is large enough.

2. Notation

The sign “ \doteq ” must be read as “equal by definition”.

Let $G=(V, E)$ be a graph and $V_0 \subset V$. Then $G(V_0)$ and $G - V_0$ denote the subgraphs of G induced by V_0 and $V - V_0$, resp. If $v \in V$ then $N_G(v) \doteq \{u \in V - \{v\} : (u, v) \in E\}$, $d_G(v) \doteq |N_G(v)|$, $\sigma(G) \doteq \max_{v \in V} d_G(v)$, $\delta(G) \doteq \min_{v \in V} d_G(v)$. For a graph $G=(V, E)$, $\bar{G}=(V, \bar{E})$ denotes the complement of G . Let $\mathcal{D}_k \doteq \{G=(V, E) : |E|/|V| \cong k\}$. Then $\eta(k) = \min_{G \in \mathcal{D}_k} \eta(G)$.

3. Results and comments

Theorem 1. $\eta(k) \cong k/270 \sqrt{\log k}$ for $k \cong 2$.

Corollary 2. $w(k) \cong k/540 \sqrt{\log k}$ for $k \cong 2$.

Corollary 3. Hadwiger’s conjecture is true for almost all graphs of n vertices.

Note. Corollary 3 was proved by P. Erdős, B. Bollobás and P. A. Catlin independently from each other.

Corollary 4. If k is large enough then Hadwiger’s conjecture is true for almost all graphs of n vertices and kn edges.

Corollary 5. $\min_{G=(V, E), |V|=n} (\eta(G) + \eta(\bar{G})) = \Omega(n/\sqrt{\log n})$.

Let $\nu(k)$ be the smallest possible Hadwiger number of a k -connected graph. Mader [5] proved that $\eta(k)/2 \cong \nu(k) \cong 2\eta(k)$. From these inequalities and Theorem 1, we obtain

Corollary 6. $\nu(k) = \Omega(k/\sqrt{\log k})$. ■

Previously, the class $\varepsilon_k = \left\{ G = (V, E) : |V| \geq k, |E| > k \cdot |V| - \binom{k+1}{2} \right\}$ and the function $\eta_1(k) = \min_{G \in \varepsilon_k} \eta(G)$ were studied instead of \mathcal{D}_k and $\eta(k)$ for different reasons (see e.g. [5], [8], [9]). As it is a more natural class, we formulate our results for \mathcal{D}_k . Actually, we prove a bit stronger statement than Theorem 1:

Theorem 1'. $\eta_1(k) \geq 1/270 \cdot (k/\sqrt{\log k})$ for $k \geq 2$.

From the proof of Theorem 1', it is easy to obtain a polynomial-time algorithm with the help of that any graph in ε_k can be contracted to the complete graph of $k/270 \sqrt{\log k}$ vertices.

We can see in the proof of Theorem 1' that the constant $1/270$ can be improved at the expense of the length of the proof. It is proved in [4] that

$$.064 \frac{k}{\sqrt{\log k}} \leq \eta_1(k) \leq 2.25 \frac{k}{\sqrt{\log k}}.$$

The structure of the proof is as follows. In section 4, we prove that it is sufficient to study the contraction of graphs containing much more edges than its complement. Some features of contraction of such graphs are studied in section 5. Using these features, a lower bound of the Hadwiger number of such graphs is given in section 6 by an auxiliary function. Theorem 1' is proved in section 7 on the basis of sections 4 and 6. Corollaries 2–5 are also proved in section 7.

4. Graphs with large number of edges

Let $k \geq 2$ and $\alpha \geq 0$ be real numbers and let

$$\mathcal{D}(k, \alpha) = \{G = (V, E) : |V| \geq 3, |E| > k \cdot |V| + \alpha \cdot |V|^2 - 1\}.$$

Lemma 1. *Suppose that $k \geq 2$ and $\alpha \geq 0$ are fixed real numbers and the graph $G = (V, E) \in \mathcal{D}(k, \alpha)$ cannot be contracted to any graph in $\mathcal{D}(k, \alpha)$. Then every edge of G is in more than $k + 2\alpha \cdot |V| - \alpha - 1$ triangles.*

Proof. Let $k \geq 2$. If $n = 3$ then $|E| \geq 2 \cdot 3 - 1 = 5$ and this is impossible. Hence $n > 3$. Suppose that an edge (u, v) of G is contained by at most $k + 2\alpha n - \alpha - 1$ triangles. Let $G' = (V', E')$ denote the graph obtained from G by the contraction of the edge (u, v) . Then $|V'| = n - 1 \geq 3$ and $|E'| \geq |E| - (k + 2\alpha n - \alpha - 1) - 1 \geq kn + \alpha n^2 - 1 - k - 2\alpha n + \alpha = k(n - 1) + \alpha(n - 1)^2 - 1$. Thus $G' \in \mathcal{D}(k, \alpha)$, a contradiction. ■

Lemma 2. *If $k \geq 39$ then every graph $G \in \mathcal{D}(k, 1/3)$ can be contracted to a graph $G_0 = (V_0, E_0)$ such that $|V_0| \geq 5k + 1$ and $\delta(G_0) \geq 7/8 \cdot |V_0|$.*

Proof. Suppose that the lemma is not true. Let $G = (V, E)$ be a counterexample with minimum number of edges and let $|V| = n$. Then $|E| \geq kn + n^2/3$. Since $nk + n^2/3 - 1 < (n^2 - n)/2$ so

$$(1) \quad n \geq 6k + 1$$

Specially, since $k \geq 39$ so

$$(2) \quad n \geq 235$$

If $\delta(G) \cong 7n/8$ then $G_0 = G$ provides the statement of the lemma, a contradiction to the choice of G . Therefore

$$(3) \quad \delta(G) < 7n/8$$

Let $v_0 \in V$, $d_G(v_0) = \delta(G)$. Consider the graph $G' \doteq G(N_G(v_0) \cup \{v_0\})$. For every vertex $v \in N_G(v_0)$, the edge (v_0, v) is in more than $k + 2n/3 - 1/3 - 1$ triangles by Lemma 1, i.e. $\delta(G') > k + 2/3 - 1/3$. Thus if $8/7(k + 2/3 - 1/3) \cong 1 + \delta(G)$ then for $G_0 = G' = (V', E')$, we have $|V'| \cong k + 2/3n - 1/3 + 1 \cong k + 4k + 2/3 + 2/3 > 5k + 1$, a contradiction to the choice of G . Thus

$$(4) \quad \frac{8}{7} \left(k + \frac{2}{3}n - \frac{1}{3} \right) < 1 + \delta(G).$$

Then

$$\frac{7}{8} \left(k + \frac{2}{3}n - \frac{1}{3} \right) < 1 + \frac{2|E|}{n} \cong 1 + 2 + \frac{2}{3}.$$

and hence

$$(5) \quad k > \frac{n}{9} - \frac{29}{18}.$$

Combining the inequalities (3), (4) and (5), we get $n < 232$, a contradiction to (2). ■

Lemma 3. *If $k \cong 8$ then every graph $G \in \mathcal{D}(k, 2/9)$ can be contracted to a graph $G_0 = (V_0, E_0)$ such that $|V_0| \cong 13k/5$ and $\delta(G_0) \cong 3|V_0|/4$.*

Proof. Suppose that the lemma is not true. Let $G = (V, E)$ be a counterexample with minimum number of edges and let $|V| = n$. As in the proof of Lemma 2, we obtain the following inequalities:

$$(1') \quad n \cong \frac{18}{5}k$$

$$(2') \quad n \cong 29$$

$$(3') \quad \delta(G) < \frac{3}{4}n$$

$$(4') \quad \frac{4}{3} \left(k + \frac{4}{9}n - \frac{2}{9} \right) < 1 + \delta(G)$$

$$(5') \quad k > \frac{2}{9}n - \frac{35}{18}.$$

Combining the inequalities (3'), (4') and (5'), we get $n < 14$, a contradiction to (2'). ■

Lemma 4. *Let $k \cong 6480$ be a natural number and let $G \in \mathcal{E}_k$ be a graph where $\mathcal{E}_k = \left\{ G = (V, E) : |V| \cong k, |E| > k \cdot |V| - \binom{k+1}{2} \right\}$. Then G can be contracted to a graph $H = (V_0, E_0)$ such that $|V_0| \cong 13k/432 + 1$ and $\delta(H) \cong 7|V_0|/8$.*

Proof. Let $G_1=(V_1, E_1)$ be a graph in \mathcal{E}_k with minimum number of edges that G can be contracted onto. Then $|E_1| < k \cdot |V_1|$ and $\delta(G_1) \leq 2k - 1$. Let $v_0 \in V_1$ be a vertex such that $\delta_{G_1}(v_0) = \delta(G_1)$ and $G_2=(V_2, E_2) \doteq G_1(N_{G_1}(v_0) \cup \{v\})$. As in the proof of Lemma 1, it is easy to see that every edge in E_1 is contained by more than $k - 1$ triangles. Thus $\delta(G_2) > k$. Notice that $|V_2| = \delta(G_1) + 1 \leq 2k$ and $|E_2| \geq 1/2\delta(G_2)|V_2| \geq \geq k|V_2|/2$.

Let $k_1 \doteq k/18$. Then $G_2 \in \mathcal{D}(k_1, 2/9)$. According to Lemma 3, G_2 can be contracted to a graph $G_3=(V_3, E_3)$ such that $|V_3| \geq 13k_1/5 = 13k/90$ and $\delta(G_3) \geq 3|V_3|/4$. Let $k_2 \doteq |V_3|/24$. Then $G_3 \in \mathcal{D}(k_2, 1/3)$ and according to Lemma 2, G_3 can be contracted to a graph $H=(V_0, E_0)$ such that $|V_0| \geq 5k_2 + 1 \geq 5/24|V_3| + 1 \geq 13k/24 \cdot 18k + 1$ and $\delta(H) \geq 7|V_0|/8$. ■

So Lemma 4 implies that to prove the inequality $\min_{G \in \mathcal{E}_k} \eta(G) \geq c_1 k / \sqrt{\log k}$ it is sufficient to show that $\eta(G) \geq c_2 n / \sqrt{\log n}$ holds for any graph $G=(V, E)$ with $|V|=4$ and $\delta(G) \geq 7n/8$.

5. Contraction of graphs with large number of edges

Lemma 5. $\eta(\bar{G}) \geq 3|V|/4 - 1/4$ holds for any graph $G=(V, E)$ with $|E| < |V|/2$.

Proof. Let $G=(V, E)$ be a counterexample with the minimum number of vertices and let $|V|=n$. It is obvious that $n > 4$. We distinguish three cases.

Case 1. $\sigma(G) \geq 2$.

Let $d_G(v_1) = \delta(G)$ and $G_1=(V_1, E_1) \doteq G - \{v_1\}$. Then $2|E_1| < n - 4| = |V_1| - 3$ and there exist isolated vertices v_2, v_3, v_4 in G_1 . Let $G_2 \doteq G_1 - \{v_1, v_3, v_4\}$. Then $\eta(\bar{G}_2) \geq (3/4)(n - 4) - 1/4 = 3n/4 - 13/4$ by the minimality of G . Since the vertices v_2, v_3, v_4 are pairwise adjacent in \bar{G}_1 so $\eta(\bar{G}) \geq \eta(\bar{G}_1) \geq 3 + \eta(\bar{G}_2) \geq 3n/4 - 1/4$, a contradiction to the choice of G .

Case 2. $|E| \leq 1$.

Then $\eta(\bar{G}) \geq n - 1 \geq 3n/4 - 1/4$.

Case 3. $\sigma(G) = 1, |E| \geq 2$.

Let $e_1=(v_1, v_2) \in E$ and $e_2=(v_3, v_4) \in E$. Contracting the edge (v_1, v_2) in the graph \bar{G} , we obtain a graph $\bar{G}_1=(V_1, \bar{E}_1)$ with $|V_1|=n - 1$ in which the vertices v_2, v_4 and the vertex \bar{v} obtained from v_1 and v_3 are adjacent to each other. Then $\eta(\bar{G} - \{v_1, v_2, v_3, v_4\}) \geq (3/4)(n - 4) - 1/4 = 3n/4 - 13/4$ by the minimality of G . Hence $\eta(\bar{G}) \geq 3 + (3n/4 - 13/4) = 3n/4 - 1/4$. ■

Lemma 6. (crucial one) Let $r \geq G$ and $G=(V, E)$ be a graph of n vertices such that $\sigma(G) \leq n/r$. Then \bar{G} can be contracted to a graph $\bar{H}=(V_0, \bar{E}_0)$ of $\lfloor n/2 \rfloor$ vertices such that $2|E_0| \leq (27r/(r - 2)^3(r - 1)^2) \cdot |V_0|^2$.

Proof. Connect every vertex $v \in V$ by a quasi-edge to the vertices which are not adjacent to v and from which more than $3n/(r - 1)(r - 2)$ paths of length two lead to v . We show that every vertex is incident to at most $(n(r - 2)/3r) - 1$ quasi-edges. If it is not the case then the number of paths of length two from v is greater than $(3n/(r - 1)(r - 2))(n(r - 2) - 3r)/3r$. At the same time, the total number of these paths

is at most $\sigma(G)(\sigma(G)-1) \leq n(n-r)/r^2$ and the inequality

$$\frac{3n}{(r-1)(r-2)} \cdot \frac{n(r-2)-3r}{3r} > \frac{n(n-r)}{r^2}$$

holds for any natural number $n \geq r \geq G$.

Let $G' = (V', E')$ denote the graph obtained from G by adding the quasi-edges. Then $\sigma(G') \leq n/r + (n(r-2)/3r) - 1 = (n(r+1/3r) - 1) < (n-1)/2$, i.e. $\overline{G'}$ has a Hamiltonian cycle and a matching of $\lfloor n/2 \rfloor$ edges.

The contraction of an edge (v_1, v_2) in \overline{G} corresponds to a sticking in G : we replace the non-adjacent vertices v_1 and v_2 with a new vertex v_3 adjacent to $w \in V - \{v_1, v_2\}$ iff $(v_1, w) \in E$ and $(v_2, w) \in E$.

Let π be an arbitrary matching of $\lfloor n/2 \rfloor$ edges in \overline{G} . Sticking every pair of vertices in G corresponding to the edges of π and deleting the vertex not covered by π if n is odd, we obtain a graph H_π of $\lfloor n/2 \rfloor$ vertices. From among the graphs H_π , we choose a graph $H \doteq H_{\pi_0} = (V_0, E_0)$ with minimum number of edges. Let $V_0 = \{v_i : i = 1, \dots, n/2\}$ and suppose that v_i is obtained by sticking v_{1i} and v_{2i} . Notice that $(v_i, v_j) \in E_0$ iff $\{(v_{1i}, v_{1j}), (v_{1i}, v_{2j}), (v_{2i}, v_{1j}), (v_{2i}, v_{2j})\} \subset E$.

A path (x_1, x_3, x_2) of length 2 will be called a fork on the vertex pair $\{x_1, x_2\}$. We will say that the pair $\{x_1, x_2\}$ is the support of this fork.

If $(v_{1i}, v_{1j}) \notin E'$ and $(v_{2i}, v_{2j}) \notin E'$ ($(v_{1i}, v_{2j}) \notin E'$ and $(v_{2i}, v_{1j}) \notin E'$, resp.) (for some $i, j \in \{1, \dots, \lfloor n/2 \rfloor\}$, $i \neq j$) then the graph $H_{ij} = (V_{ij}^1, E_{ij}^1)$ ($H_{ij}^2 = (V_{ij}^2, E_{ij}^2)$, resp.) is called the $(i, j, 1)$ -transformation ($(i, j, 2)$ -transformation, resp.) of the graph $H_0 = (E_0, V_0)$. The graph H_{ij}^1 (H_{ij}^2 , resp.) is obtained by stickings along the matching π_{ij}^1 (π_{ij}^2 , resp.) in $\overline{G'}$, where π_{ij}^1 (π_{ij}^2 , resp.) is obtained from π_0 replacing the edges (v_{1i}, v_{2i}) and (v_{1j}, v_{2j}) with the edges (v_{1i}, v_{1j}) and (v_{2i}, v_{2j}) ((v_{1i}, v_{2j}) and (v_{2i}, v_{1j}) , resp.). These transformations will be called (i, j) -transformations.

Let $d_H(v_i) = t_i$ ($1 \leq i \leq \lfloor n/2 \rfloor$). During an (i, j, l) -transformation, $t_i + t_j$ edges of $E_0 - E_{ij}^l$ are destroyed but it may occur that new edges arise. The edges of G' hinder the existence of some (i, j) -transformations. Since every edge of G' may cross at most one transformation so every vertex $v_i \in V$ can take part in at least $n - 3 - 2 \left(\frac{n(r+1)}{3r} - 1 \right)$ transformations. Then $|E_{ij}^l - E_0| \geq t_i + t_j$ holds for $1 \leq i, j \leq$

$\lfloor n/2 \rfloor$, $1 \leq l \leq 2$ by the choice of H . Every edge of $E_{ij}^l - E_0$ has a pair of forks in G with common supports in $\pi_0 \cap \pi_{ij}^l$. Every pair of forks may take part in at most one transformation producing "new" edge. Then every (i, j, l) -transformation corresponds to $x_{ij}^l \geq t_i + t_j$ pairs of forks and the members of every such pair have a common support belonging to π_0 . Associate t_i pairs from among these x_{ij}^l pairs of forks with v_i and t_j pairs from among the remaining $x_{ij}^l - t_i$ pairs with v_j . Then at least $t_i(n - (2n(r+1)/3r) - 1)$ pairs of forks will belong to v_i ($1 \leq i \leq \lfloor n/2 \rfloor$). Furthermore, every edge $(v_i, v_j) \in E_0$ corresponds to the pairs $\{(v_{1i}, v_{1j}, v_{2i}, v_{2j}, v_{2i})\}$ and $\{(v_{1j}, v_{1i}, v_{2j}, v_{2i}, v_{2j})\}$ of forks not belonging to any vertex. I.e., the total number of pairs of forks in G with common supports in π_0 is at least

$$\sum_{i=1}^{\lfloor \frac{n}{2} \rfloor} t_i \left(\frac{n(r-2)}{3r} - 1 \right) + 2|E_0| = \left(\frac{n(r-2)}{3r} - 1 \right) \cdot 2|E_0| + 2|E_0| = 2|E_0| \frac{n(r-2)}{3r}.$$

On the other hand, $(v_{1i}, v_{2i}) \in E'$ for $1 \leq i \leq \lfloor n/2 \rfloor$ thus $\{v_{1i}, v_{2i}\}$ is the support of at most $(3n/(r-1)(r-2))$ forks and so it is the common support of at most $\binom{3n}{(r-1)(r-2)}$ pairs of forks. Hence

$$\frac{2|E_0|n(r-2)}{3r} \leq \left\lfloor \frac{n}{2} \right\rfloor \frac{9n}{2(r-1)^2(r-2)^2} \cdot \left(n - \frac{(r-1)(r-2)}{3} \right), \text{ i.e.}$$

$$2|E_0| \leq \left\lfloor \frac{n}{2} \right\rfloor \cdot \frac{n-1}{2} \cdot \frac{27r}{(r-1)^2(r-2)^3} \cdot \blacksquare$$

6. Lower estimation of $\eta(\bar{G})$

Let $f(n, m) \doteq n/2\sqrt{\log n}(\sqrt{\log n^2/2m} - 1.4)$ and $f(G) \doteq f(|V|, |E|)$ for any graph $G=(V, E)$. The aim of this section is to give a lower bound of $\eta(\bar{G})$ by $f(G)$ for graphs with small ratio $|E|/|V|^2$.

Lemma 7. Let $n \geq 3, 2m \geq n, 1.4 < \sqrt{\log n^2/2m}$ and $n/2 > \alpha \geq 1 + \log n^2/2m - 1.4\sqrt{\log n^2/2m}$. Then $f(n, m) \leq f(n-1, m-2m\alpha/n)$.

Proof. Let $c \doteq n^2/2m, c_1 \doteq (n-1)^2/2(m-2m\alpha/n)$. Then

$$f\left(n-1, m-\frac{2m\alpha}{n}\right) - f(n, m) \geq \frac{1}{2\sqrt{\log n}} \left((n-1)(\sqrt{\log c_1} - 1.4) - n(\sqrt{\log c} - 1.4) \right).$$

I.e., it is sufficient to prove the inequality

$$(6) \quad (n-1)\sqrt{\log c_1} \geq n\sqrt{\log c} - 1.4.$$

The inequality for α in the conditions of the lemma implies that

$$2 \log c - 2.8\sqrt{\log c} \geq 2(\alpha-1) \cdot \frac{(n-1)^2}{n(n-2)}.$$

Since $x < \log(1-x)^{-1}$ for $x \in (0, 1)$ so

$$\begin{aligned} 2(\alpha-1) \frac{(n-1)^2}{n(n+2)} &\leq \frac{(n-1)^2}{n} \log \left(1 - \frac{2(\alpha-1)}{n-2} \right)^{-1} = \frac{(n-1)^2}{n} \cdot \log \frac{n-2}{n-2\alpha} \leq \\ &\leq \frac{(n-1)^2}{n} \cdot \log \frac{\left(1 - \frac{1}{n}\right)^2}{1 - \frac{2\alpha}{n}} = \frac{(n-1)^2}{n} \cdot \left[\log \frac{n^2 \left(1 - \frac{1}{n}\right)^2}{2m \left(1 - \frac{2\alpha}{n}\right)} - \log \frac{n^2}{2m} \right] = \\ &= \frac{(n-1)^2}{n} (\log c_1 - \log c). \end{aligned}$$

Thus

$$(n-1)^2 (\log c_1 - \log c) \geq 2n \log c - 2.8n \sqrt{\log c}.$$

Then

$$\begin{aligned} (n-1)^2 \log c_1 &\cong \\ &\cong (n-1)^2 \log c + 2n \log c - 2.8n \sqrt{\log c} = n^2 \log c + \log c - 2.8n \sqrt{\log c} = \\ &= (n \sqrt{\log c} - 1.4)^2 + (\log c - 1.96) > (n \sqrt{\log c} - 1.4)^2, \end{aligned}$$

i.e. (6) holds. ■

With the help of derivates, it is easy to prove

Lemma 8. *If $x \cong \sqrt{\log 49}$ then $e^{1.4x}/5 \cong 1 + x^2 - 1.4x$. ■*

Lemma 9. *Let $G=(V, E)$ be a graph with $|V|=n, E \neq \emptyset$ and suppose that $c \doteq n^2/2|E| \cong \cong 49, \sigma(G) \cong n/c \cdot (1/5)e^{1.4\sqrt{\log c}}$. Then \bar{G} can be contracted to a graph $\bar{H}=(v_0, \bar{E}_0)$ of $\lfloor n/2 \rfloor$ vertices such that either $E_0 = \emptyset$ or $f(H) \cong f(G)$.*

Proof. Let $z \doteq \sigma(G)c/n$. According to Lemma 6, \bar{G} can be contracted to a graph $\bar{H}=(V_0, \bar{G}_0)$ such that $|V_0|=\lfloor n/2 \rfloor$ and either $E_0 = \emptyset$ or $c_1 \doteq |V_0|^2/2|E_0| \cong (c/z)^4 \cdot \cdot 1/27 \cdot (1-3/c)^2(1-23/c)^3$. Suppose that $E_0 \neq \emptyset$. If $c \cong 49$ then $(1/5c)e^{1.4\sqrt{\log c}} < 1/15$ and if $z/c < 1/15$ holds then $(1-z/c)^2(1-2z/c)^3 > 43/76$. Thus

$$(7) \quad c_1 \cong \frac{43}{76 \cdot 27} \cdot \left(\frac{c}{z}\right)^4 > 0.02 \left(\frac{c}{z}\right)^4.$$

Estimate $f(H) - f(G)$.

$$\begin{aligned} f(H) - f(G) &= \frac{\lfloor \frac{n}{2} \rfloor (\sqrt{\log c_1} - 1.4)}{2 \sqrt{\log \lfloor \frac{n}{2} \rfloor}} - \frac{n(\sqrt{\log c} - 1.4)}{2 \sqrt{\log c}} \cong \\ &\cong \frac{n}{2 \sqrt{\log n}} \left(\frac{1}{2} \sqrt{\log c_1} - \sqrt{\log c} + 0.7 \right) = \frac{n}{2 \sqrt{\log n}} \frac{0.25 \log c_1 - (\sqrt{\log c} - 0.7)^2}{0.5 \sqrt{\log c_1} + \sqrt{\log c} - 0.7}. \end{aligned}$$

I.e., it is sufficient to prove the inequality

$$(8) \quad \frac{1}{4} \log c_1 - \log c + 1.4 \sqrt{\log c} - .49 \cong 0.$$

Applying (7), the left side of (8) is at least

$$\log c - \log z + \frac{1}{4} \log .02 - \log c + 1.4 \sqrt{\log c} - 0.49 \cong 1.4 \sqrt{\log c} - \log(z \cdot 50^{1/4} e^{1/2}).$$

We have $z \cong e^{1.4\sqrt{\log c}}/5 < e^{1.4\sqrt{\log c}}/\sqrt[4]{50e^2}$ by the conditions of the lemma and so (8) holds. ■

Lemma 10. *The inequality $\eta(\bar{G}) \cong f(G)$ holds for any graph with $E \neq \emptyset$ and $|V|^2/2|E| \cong \cong 49$.*

Proof. Let $|V|=n, c \leq |V|^2/2|E|, z = \sigma(G) \cdot c/n$. If $n \leq 8$ then the statement of the lemma is obvious. Suppose that the lemma is true for $n' < n$. We distinguish four cases.

Case 1. $0 < 2|E| < n$.

If $|E| > 0$ then $f(G) \leq n\sqrt{\log n^2/2\sqrt{\log n}} = (\sqrt{2}/2)n$ and $\eta(\bar{G}) \leq 3n/4 - 1/4 \leq (\sqrt{2}/2)n$ by Lemma 5.

Case 2. $2|E| \leq n, z \leq e^{1.4\sqrt{\log c}}/5$.

Since $c \leq n$ so $f(G) \leq |n/2|$. Let $H = (V_0, E_0)$ be a graph described in Lemma 9. If $E_0 = \emptyset$ then $\eta(\bar{G}) \leq |n/2| \leq f(G)$. If $E_0 \neq \emptyset$ then $f(H) \leq f(G)$ and the characteristics of the function f imply that $(|V_0|^2/2|E_0|) > c \geq 49$. Then $\eta(\bar{H}) \leq f(H) \leq f(G)$ by the induction hypothesis. But $\eta(\bar{G}) \leq \eta(\bar{H})$.

Case 3. $2|E| \leq n, z \geq 1 + \log c - 1.4\sqrt{\log c}$.

Deleting a vertex v_0 with $d_G(v_0) = \sigma(G) = zn/c$ from G , we obtain a graph $G' = (V', E')$ with $|V'| = n - 1$ and $|E'| \leq |E| - 2|E|/n(1 + \log c - 1.4\sqrt{\log c})$. Now $f(G') \leq f(G)$ by Lemma 7. It is clear that $|V'|^2/2|E'| \geq c \geq 49$. And $\eta(\bar{G}') \leq f(G') \leq f(G)$ by the induction hypothesis.

Thus, we have to settle only

Case 4. $2|E| \leq n, 1 + \log c - 1.4\sqrt{\log c} > z > e^{1.4\sqrt{\log c}}/5$.

If $c \geq 49$ then $e^{1.4\sqrt{\log c}}/5 \geq 1 + \log c - 1.4\sqrt{\log c}$ holds by Lemma 8. ■

7. Proof of Theorem 1' and its corollaries

Proof of Theorem 1'.

Case 1. $k \leq 6480$.

Mader [5] proved that $\eta_1(k) \leq 1/8(k/\log_2 k)$ and if $k \leq 6480$ then $(1/8)(k/\log k) \leq (1/270)(k/\sqrt{\log k})$ holds.

Case 2. $k > 6480$.

Let $G \in \mathcal{E}_k$. Then G can be contracted to a graph $H = (V_0, E_0)$ such that $|V_0| \leq 13k/432 + 1$ and $\sigma(\bar{H}) \leq |V_0|/8$ by Lemma 4. Now H can be contracted to a graph $H_1 = (V_1, E_1)$ such that $|V_1| \leq 13k/864$ and $2|\bar{E}_1| \leq |V_1|^2/49$ by Lemma 6. Let $c \leq (|V_1|^2/2|E_1|)$ if $\bar{E}_1 \neq \emptyset$. Then

$$\eta(H_1) \leq f(\bar{H}_1) = \frac{(\sqrt{\log c} - 1.4)|v_1|}{2\sqrt{\log |v_2|}} \leq \frac{.5|v_1|}{2\sqrt{\log |v_1|}} \leq \frac{1}{4} \cdot \frac{13k}{864} \cdot \frac{1}{\sqrt{\log k}} \leq \frac{k}{270\sqrt{\log k}}$$

by Lemma 10. Finally if $\bar{E}_1 = \emptyset$ then $\eta(G) \leq |V_1| \leq 13/864k$. ■

Proof of Corollary 2. If $k \leq 1000$ then the statement holds obviously. Let $k > 1000$. Every k -chromatic graph G contains a critical k -chromatic subgraph G' such that the

degree of every vertex of G' is at least $k-1$. It is clear that $G' \in \mathcal{D}_{\lfloor \frac{k-1}{2} \rfloor}$. Then $\eta(G) \cong$

$$\cong \eta(G') \cong 1/270 \cdot \frac{\lfloor \frac{k-1}{2} \rfloor}{\sqrt{\log k/2}} \cong k/540 \sqrt{\log k} \text{ by Theorem 1. } \blacksquare$$

Proof of Corollary 3. It is known that $\chi(G) \cong n/\log n$ holds for almost all graphs G with n vertices (see [3]). Furthermore, it is easy to see that $|E| \cong n^2/\delta$ for almost all graphs $G=(V, E)$ with n vertices. Then $\eta(G) \cong 1/270 \cdot (n/6)/\sqrt{\log n}$ by Theorem 1. If, in addition $\sqrt{\log n} > 6 \cdot 270$ then $\chi(G) \cong \eta(G)$. \blacksquare

Proof of Corollary 4. It is known (see [1], p. 96) that if k is large enough then $\chi(G) \cong 3k/\log k$ holds for almost all graphs G of n vertices and kn edges. Now $\eta(G) \cong (1/270)(k/\sqrt{\log k})$ for any graph G of this class by Theorem 1. Therefore if k is large enough then Corollary 4 is true. \blacksquare

Proof of Corollary 5. Let $G=(V, E)$ be an arbitrary graph with $|V|=n$. Then either $|E| \cong (n-1)n/4$ or $|\bar{E}| \cong (n-1)n/4$. Now $\max\{\eta(G), \eta(\bar{G})\} \cong (1/270)(n/4\sqrt{\log n})$ by Theorem 1.

On the other hand, as we said it in the introduction, the Hadwiger number of almost all graphs of n vertices is at most $n/\sqrt{\log n}$. Thus if n is large then there exists a graph $G_0=(V_0, E_0)$ with $|V_0|=n$ such that $\eta(G_0) \cong n/\sqrt{\log n}$ and $\eta(\bar{G}_0) \cong n/\sqrt{\log n}$. Then $\eta(G_0) + \eta(\bar{G}_0) \cong 2n/\sqrt{\log n}$. \blacksquare

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