Graph coloring. Part 2

Lecture 33

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Proposition 5.3. Definitions (A) and (B) are equivalent.

Proof: In the lecture.

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Examples: (a) Complete graphs, (b) Odd cycles, (c) Odd wheels, (d) Moser spindle.

Theorem 5.6. Let $k \geq 3$ and G be a k-critical graph. Then
(a) $\kappa(G) \geq 2$;
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Proof of (a): Suppose our k-critical G is disconnected and G_1 is a component of G. Since G is k-critical, its proper subgraphs G_1 and $G - V(G_1)$ have chromatic number at most $k - 1$. Let f_1 and $f - 2$ be their $(k-1)$-colorings. Then $f_1 \cup f_2$ is a $(k-1)$-coloring of G , a contradiction.

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Suppose now G is connected and v is a cut vertex in G. Let G' be a component of $G - v$, $G_1 = G - V(G')$ and $G_2 = G[V(G') + v]$. Again, since G is k-critical, for $i = 1, 2$, G_i has a $(k-1)$-coloring f_i. We can rename the colors in f_2 to make $f_2(v) = f_1(v)$. Then again, $f_1 \cup f_2$ is a $(k-1)$-coloring of G , a contradiction.

Proof of (b): Suppose $\kappa'(G) \leq k-2$. Then G has a vertex partition $(S, \overline{S})$ s.t. $E_G(S, \overline{S}) = \{x_i y_i : 1 \leq i \leq s\}$, where $s \leq k-2$, $\{x_1, \ldots, x_s\} \subset S$ and $\{y_1, \ldots, y_s\} \subset \overline{S}$.
Note that x_is do not need to be all distinct and y_is do not need to be all distinct.

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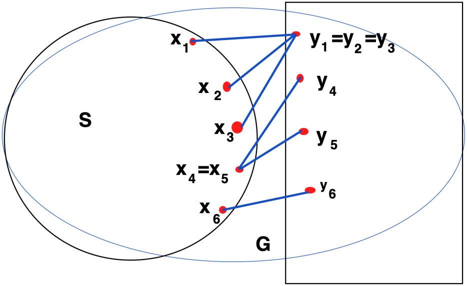
Let $G_1 = G[S]$ and $G_2 = G[\overline{S}]$. Since G is k-critical, for $i = 1, 2$, G_i has a $(k-1)$-coloring f_i.



We try to rename the colors of f_2 so that $f_1(x_i) \neq f_2(y_i)$ for all $1 \leq i \leq s$. There are $(k-1)!$ ways to rename these $k-1$ colors. Each of the edges $x_i y_i$ spoils the $(k-2)!$ cases where $f_1(x_i) = f_2(y_i)$.
Then the number of ways to rename the colors which are not spoiled is at least

$$(k-1)! - s((k-2)!) = (k-2)!((k-1)-s) \geq (k-2)! > 0.$$

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Proof of the "Moreover" part: We describe the Hájos Construction that creates from two k-critical graphs a new k-critical graph with connectivity exactly 2.

1) Take two disjoint k-critical graphs G_1 and G_2.
2) Choose an edge $x_1 y_1$ in G_1 and an edge $x_2 y_2$ in G_2.
3) Delete the edges $x_1 y_1$ and $x_2 y_2$, glue x_2 with x_1 into a new vertex x^*, add edge $y_1 y_2$. Call new graph G^*.

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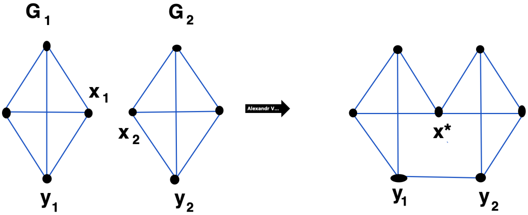

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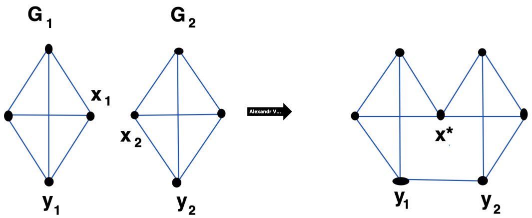

By construction, set $\{x^*, y_1\}$ is separating in G^*. So $\kappa(G^*) = 2$.

Now we show that G^* is k-critical.
Suppose G^* has a $(k-1)$-coloring f. Since $f|_{V(G_1)}$ is NOT a $(k-1)$-coloring of G_1, $f(x^*) = f(y_1)$. Similarly, $f(x^*) = f(y_2)$. But then $f(y_1) = f(y_2)$, a contradiction.

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Consider $G^* - y_1y_2$. Since G_1 and G_2 are k-critical, for $i = 1, 2$, $G_i - x_iy_i$ has a $(k-1)$-coloring f, and $f_i(y_i) = f_i(x_i)$. Then after permuting the colors in f_2 so that $f_2(x_2) = f_1(x_1)$, we get that $f_1 \cup f_2$ is a $(k-1)$-coloring of $G^* - y_1y_2$.

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Finally, let uv be any other edge of G^*. By symmetry, we may assume $\{u, v\} \subset V(G_1)$ (or one of them is x^*). Then $G_1 - uv$ has a $(k-1)$-coloring f_1 and $G_2 - x_2y_2$ has a $(k-1)$-coloring f_2. After permuting the colors in f_2 so that $f_2(x_2) = f_1(x_1)$, we get that $f_1 \cup f_2$ is a $(k-1)$-coloring of $G^* - uv$.

Definitions

Mycielski's Construction: $M_3 = C_5$. Suppose M_k is a triangle-free graph with $\chi(M_k) = k$ and $V(M_k) = V_k = \{v_1, ..., v_{n_k}\}$. Let $V_k' = \{u_1, ..., u_{n_k}\}$. Then $V(M_{k+1}) = V_k \cup V_k' \cup \{w\}$, $M_{k+1}[V_k] = M_k$, $N_{M_{k+1}}(w) = V_k'$ and for each $1 \leq j \leq n_k$, $N_{M_{k+1}}(u_j) = N_{M_k}(v_j) \cup \{w\}$.

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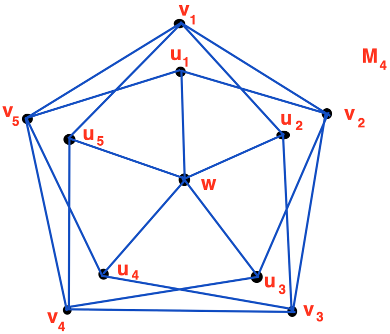

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Then color k_0 is not used on V'_{k_0}. Let $W = \{v_{i_1}, \ldots, v_{i_s}\}$ be the set of the vertices in V_{k_0} colored with k_0. We will recolor them: for each $1 \leq j \leq s$, recolor v_{i_j} with $f(u_{i_j})$. Then color k_0 is not used in the new coloring f' of M_k.

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We claim that

$$f'(v_i) \neq f'(v_j) \quad \text{for each edge } v_i v_j \in E(M_{k_0}). \qquad (1)$$

Indeed, suppose $f'(v_i) = f'(v_j)$. If $f(v_i) \neq k_0$ and $f(v_j) \neq k_0$, then the colors of v_i and v_j did not change, but $f(v_i) \neq f(v_j)$, a contradiction. If $f(v_i) = k_0 = f(v_j)$, then v_i and v_j cannot be adjacent.

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This shows that the difference $\chi(G) - \omega(G)$ and the ratio $\frac{\chi(G)}{\omega(G)}$ can be arbitrarily large.